

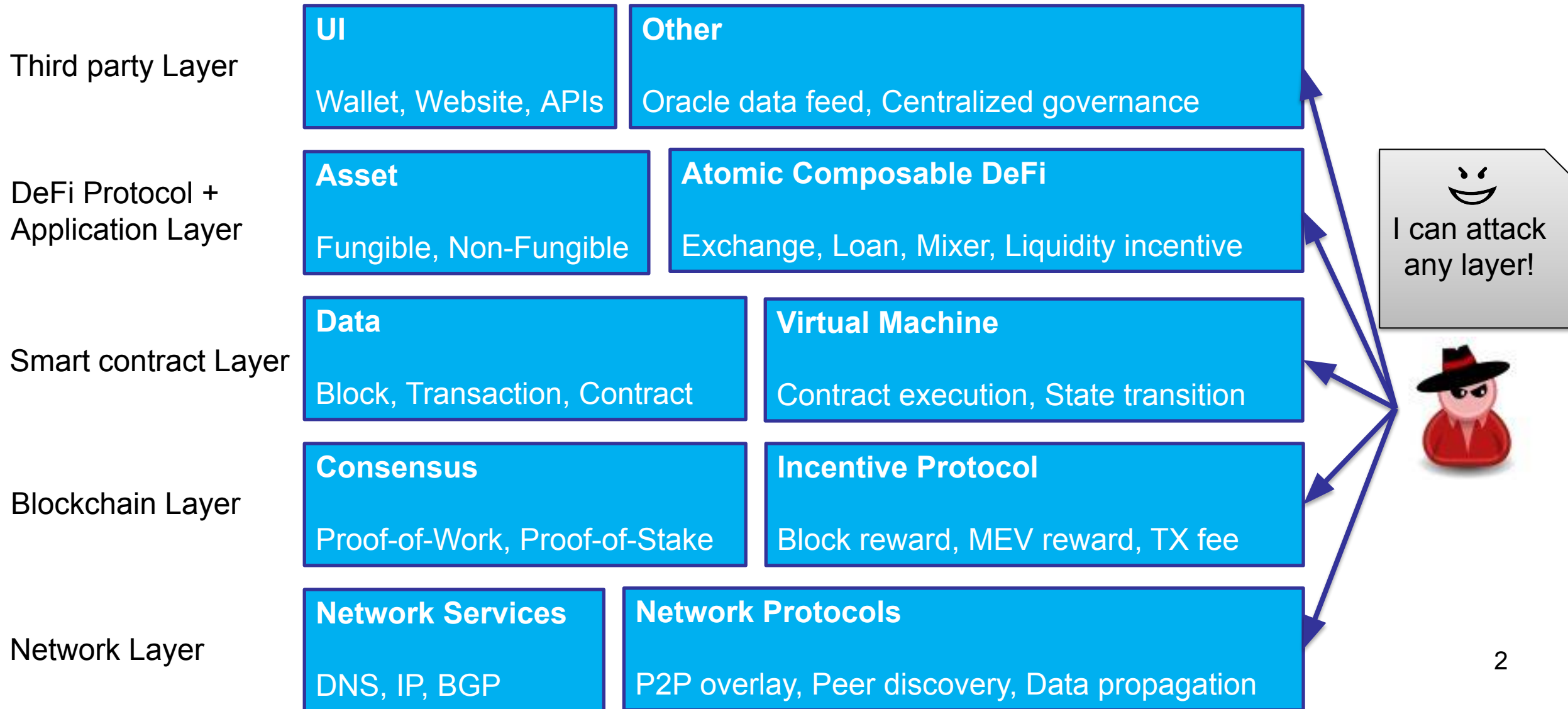
# Decentralized Finance

## DeFi Security

Instructors: Dan Boneh, Arthur Gervais, Andrew Miller, Christine Parlour, Dawn Song



# DeFi Security Affects Multiple Layer





# Network Layer Security

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# Network Layer

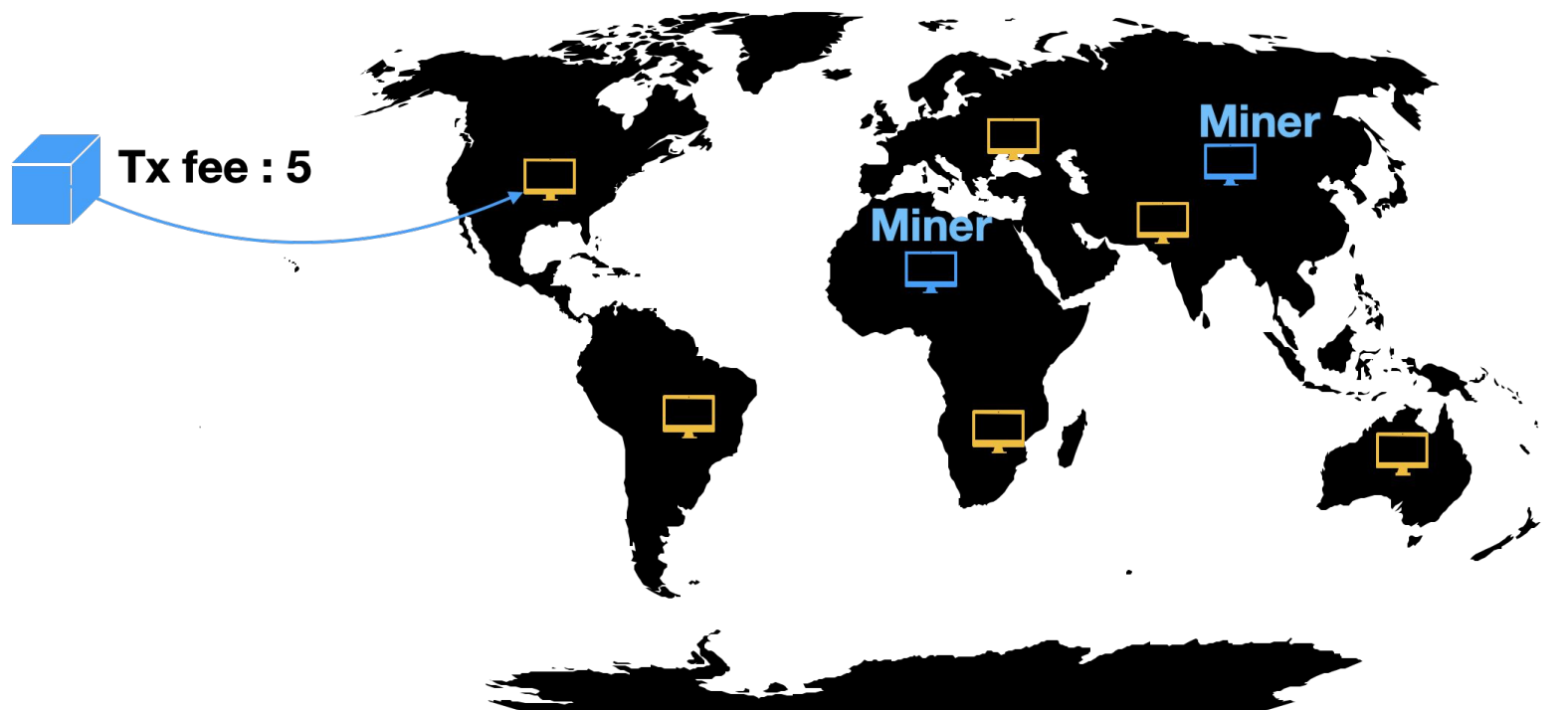
---

- Why Network Layer?
  - Information dissemination and propagation.
  - Latency matters!
- How many nodes?
  - Bitcoin: about 10'000 reachable full nodes (TCP/8333)
  - Ethereum:
  - Dogecoin:
- What type of nodes exist?
  - Full nodes
  - Light nodes

# Exchange Transaction Propagation

Trader

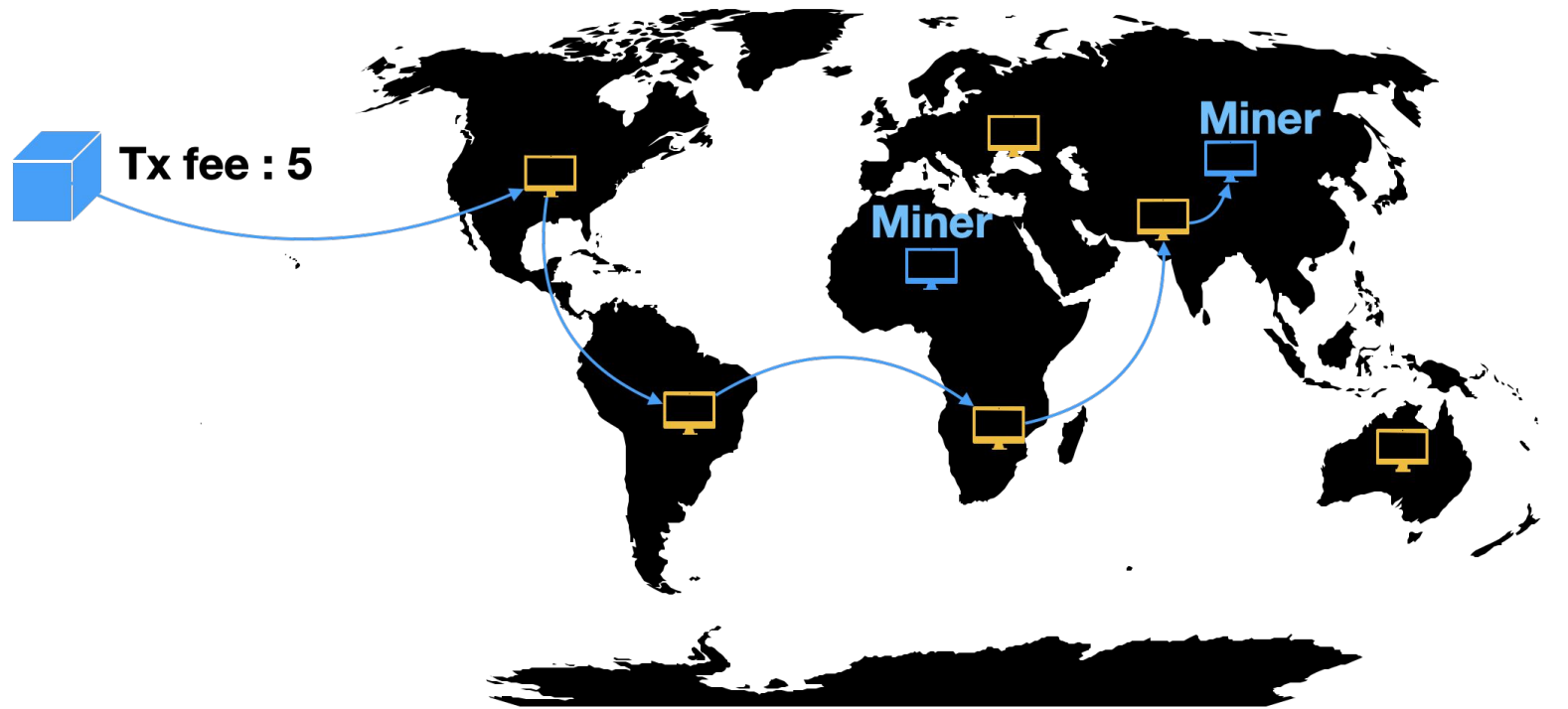
P2P Network



# Exchange Transaction Propagation

Trader

P2P Network

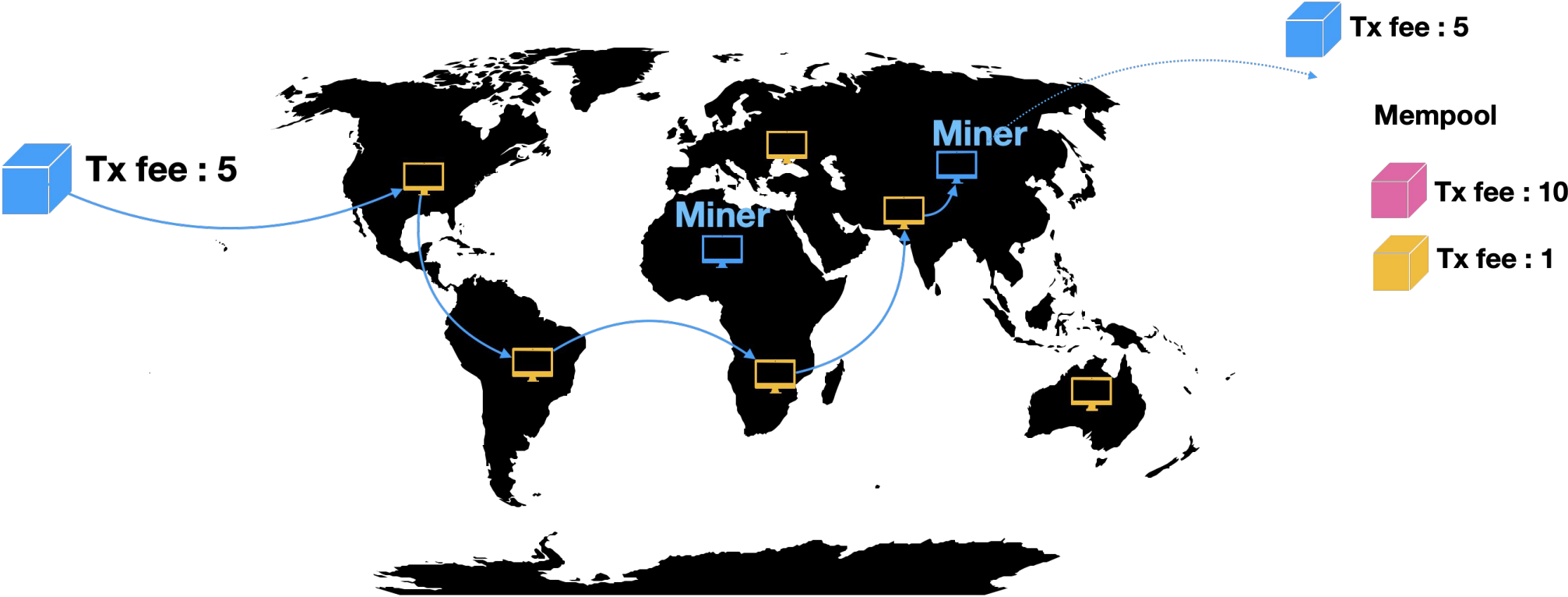


# Exchange Transaction Propagation

Trader

P2P Network

Elected Leader/Miner

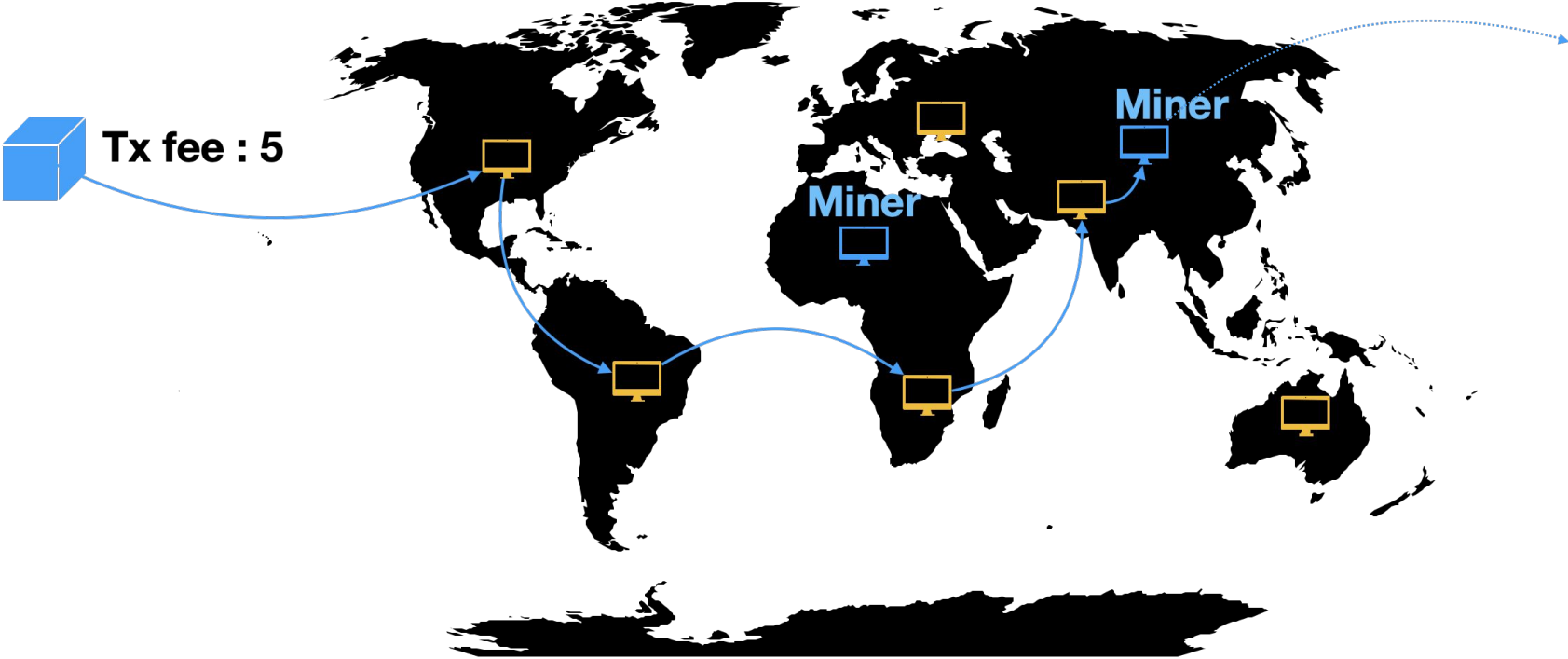


# Exchange Transaction Propagation




Trader

P2P Network

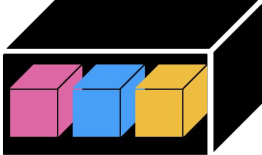
Elected Leader/Miner



Mempool

-  Tx fee : 10
-  Tx fee : 5
-  Tx fee : 1

Final Block

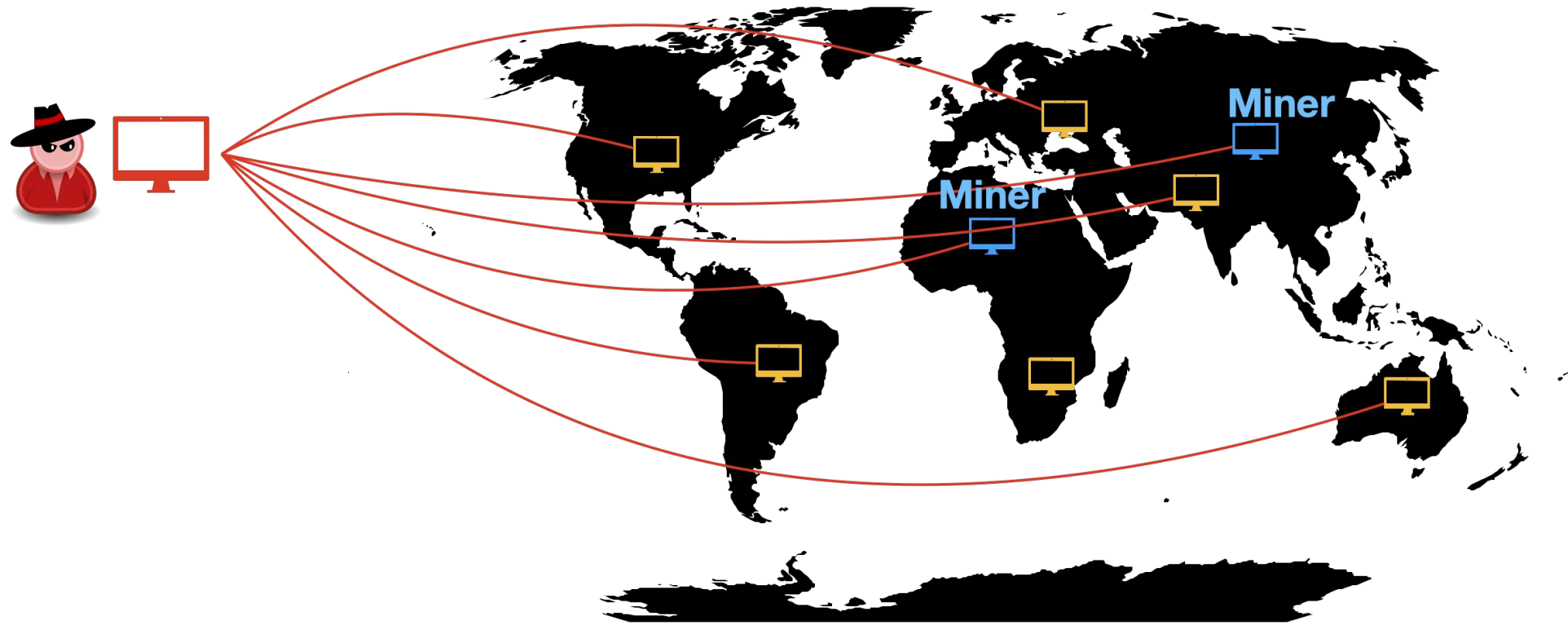




# Network Layer – Spy Node

Attacker (Spy Node)

P2P Network

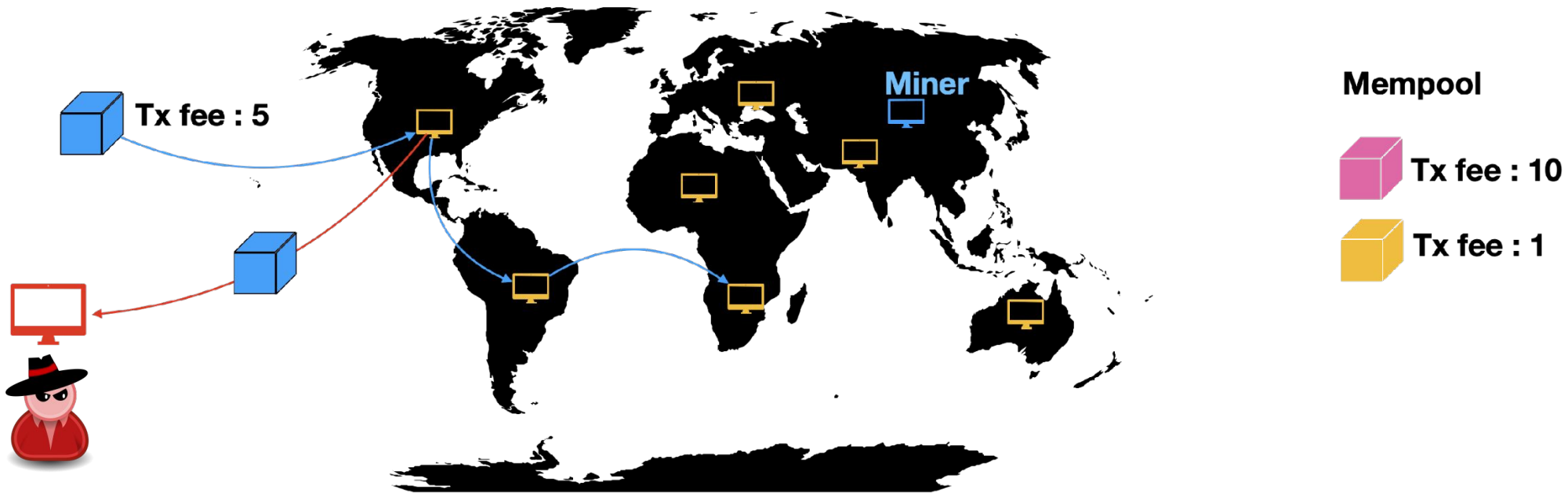


# Network Layer – Spy Node

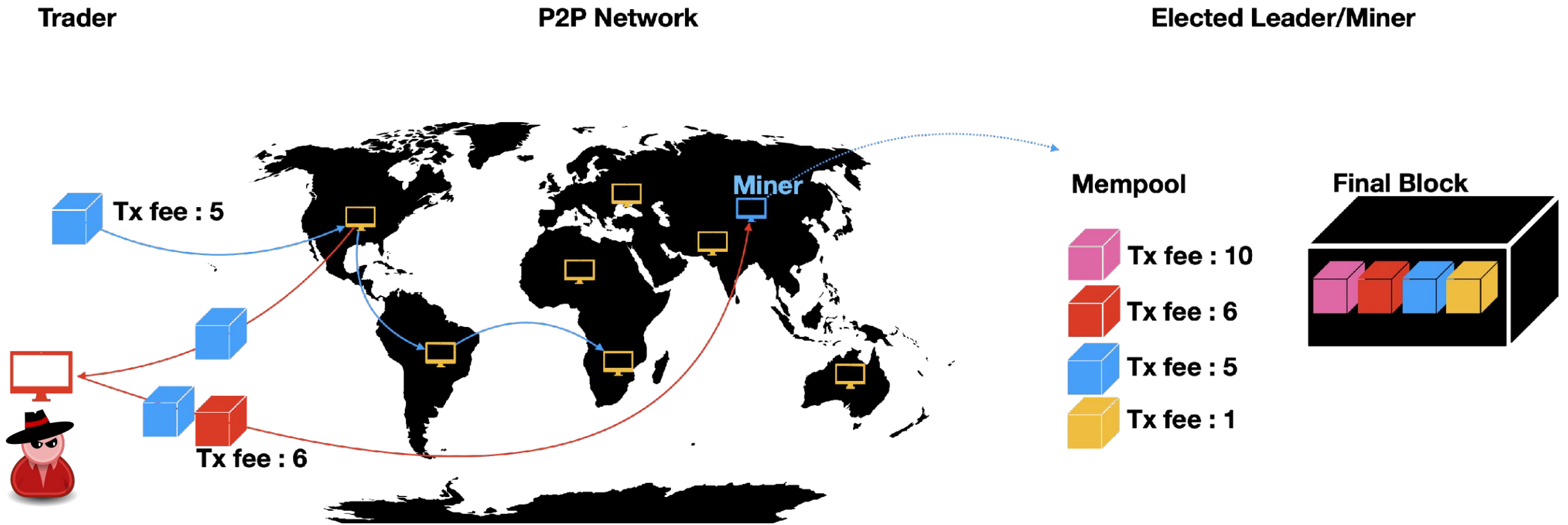
Trader

P2P Network

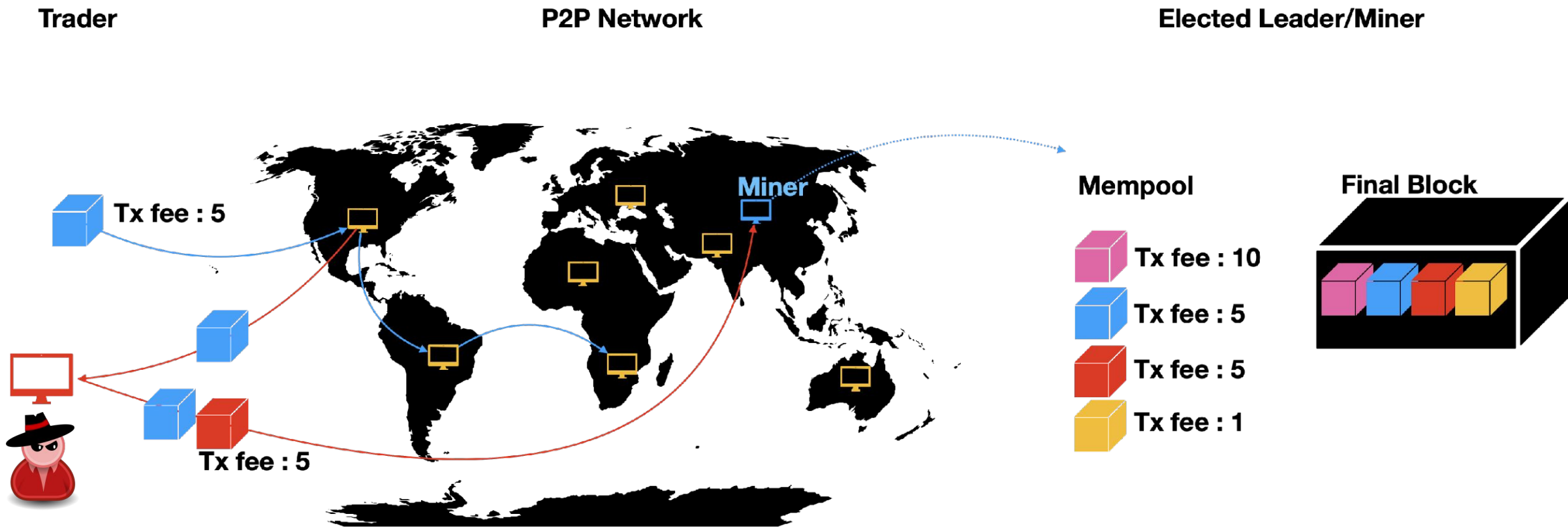
Elected Leader/Miner



# Front-running



# Back-running

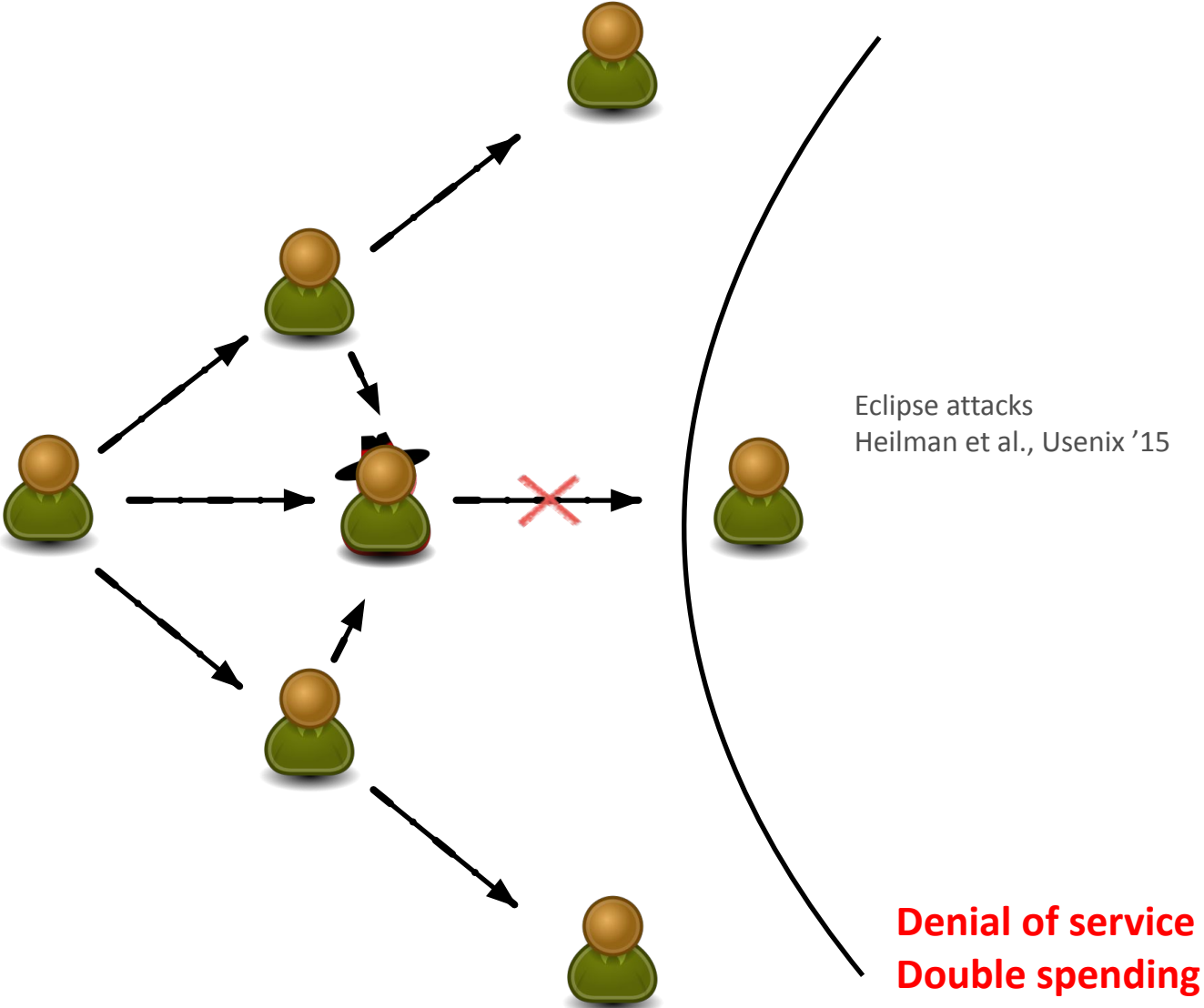




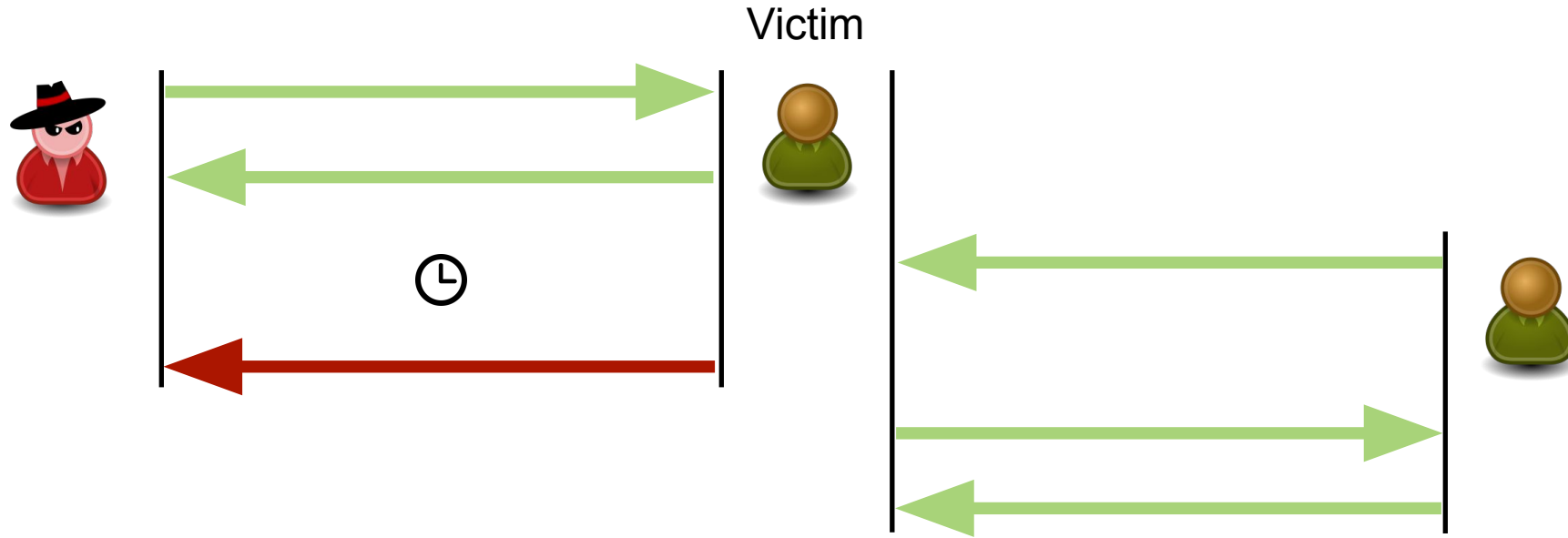
# Eclipse Attacks

---

# Eclipse Attacks



# Request timeouts



Block timeout: 20 minutes  
Transaction timeout: 2 minutes

# Security Implications

- Adversary

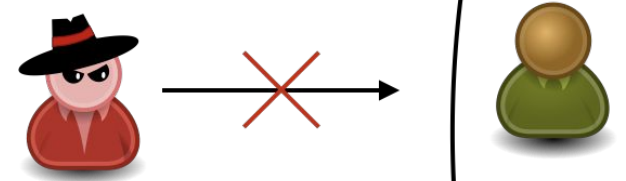
- Blinds victim from blocks and transaction > 20 min
- Experimental validation

- Impact

- Double spend transactions
- Aggravated selfish mining
- Network wide Denial of Service

- Mitigations

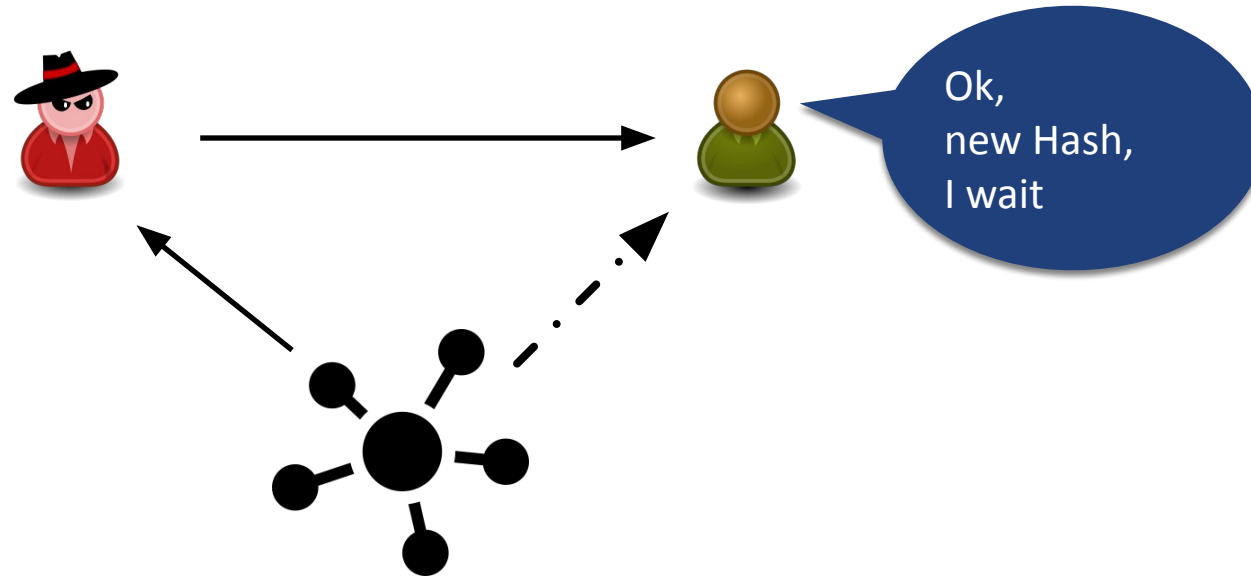
- Hardening measures
- Estimate waiting time for secure transactions





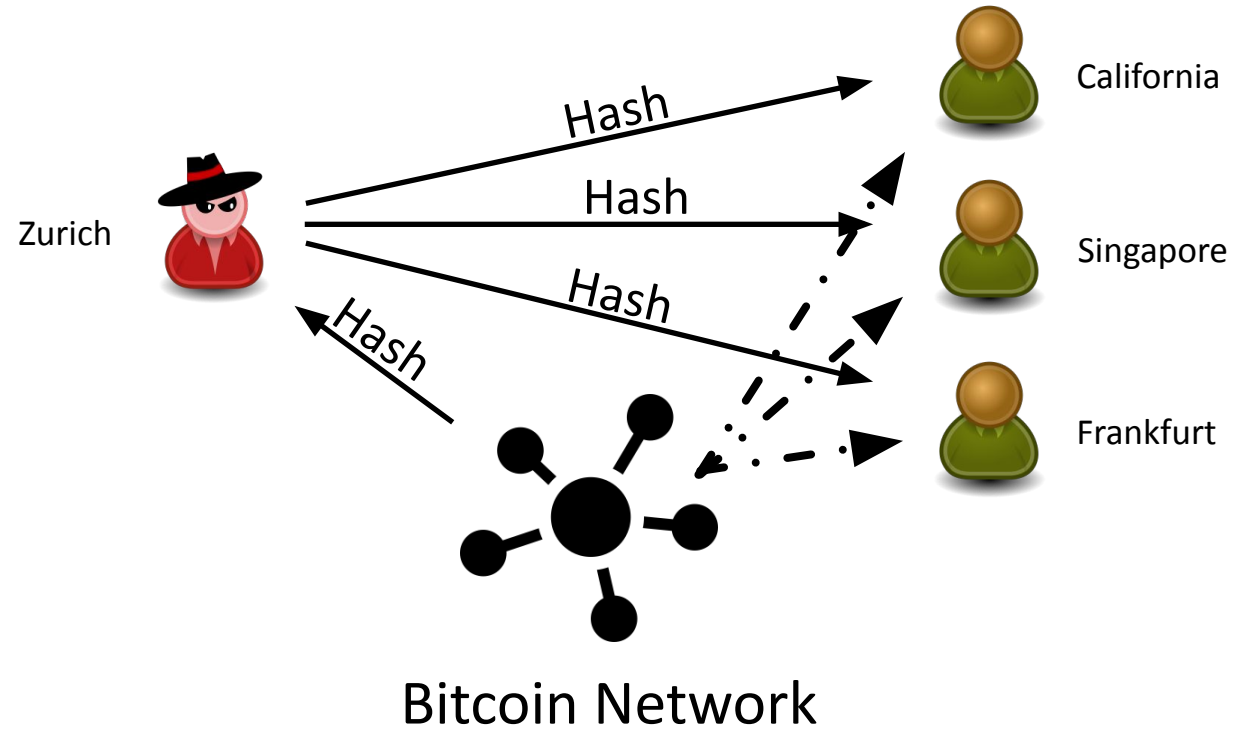
# Eclipse Requirements

1. Must be **first** peer to advertise Transaction/Block



2. Victim should wait
  - Block timeout: 20 minutes
  - Transaction timeout: 2 minutes

# Being First on the Network Layer

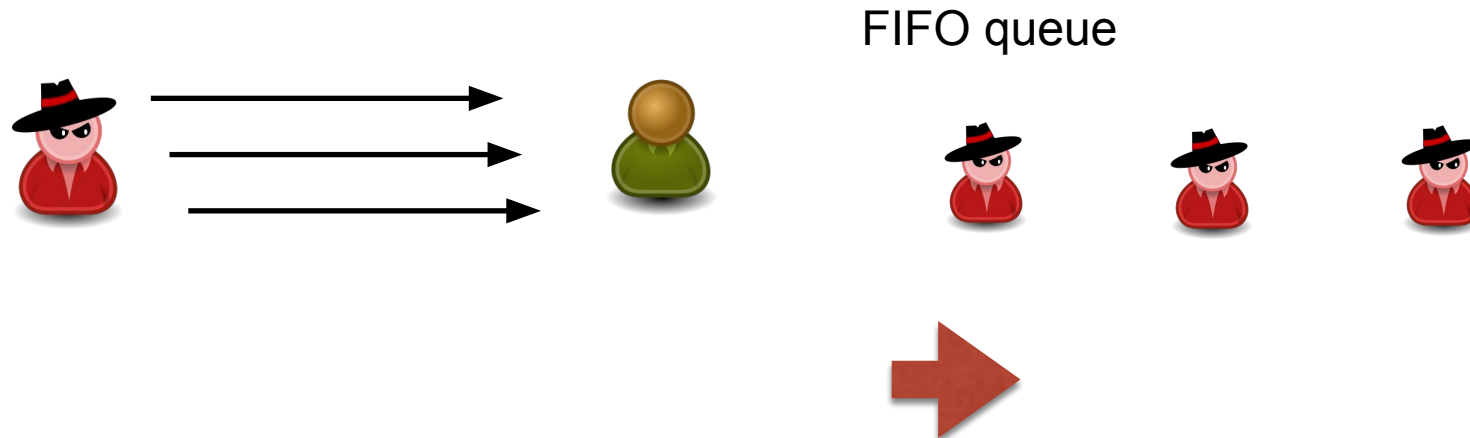


Connections of Adversary	40	80	200	800
Connections of Victim	40	40	40	40
Average success in being first	0.44± 0.14	0.57± 0.20	0.80± 0.14	<b>0.89±</b> <b>0.07</b>

# Network Layer Timeouts

- Transactions

- After 2 minutes request from other peer (FIFO)



- Blocks (older Bitcoin version)

- After 20 minutes disconnect and do nothing
- If received header, disconnect and request block from another peer



# Blockchain Layer Security

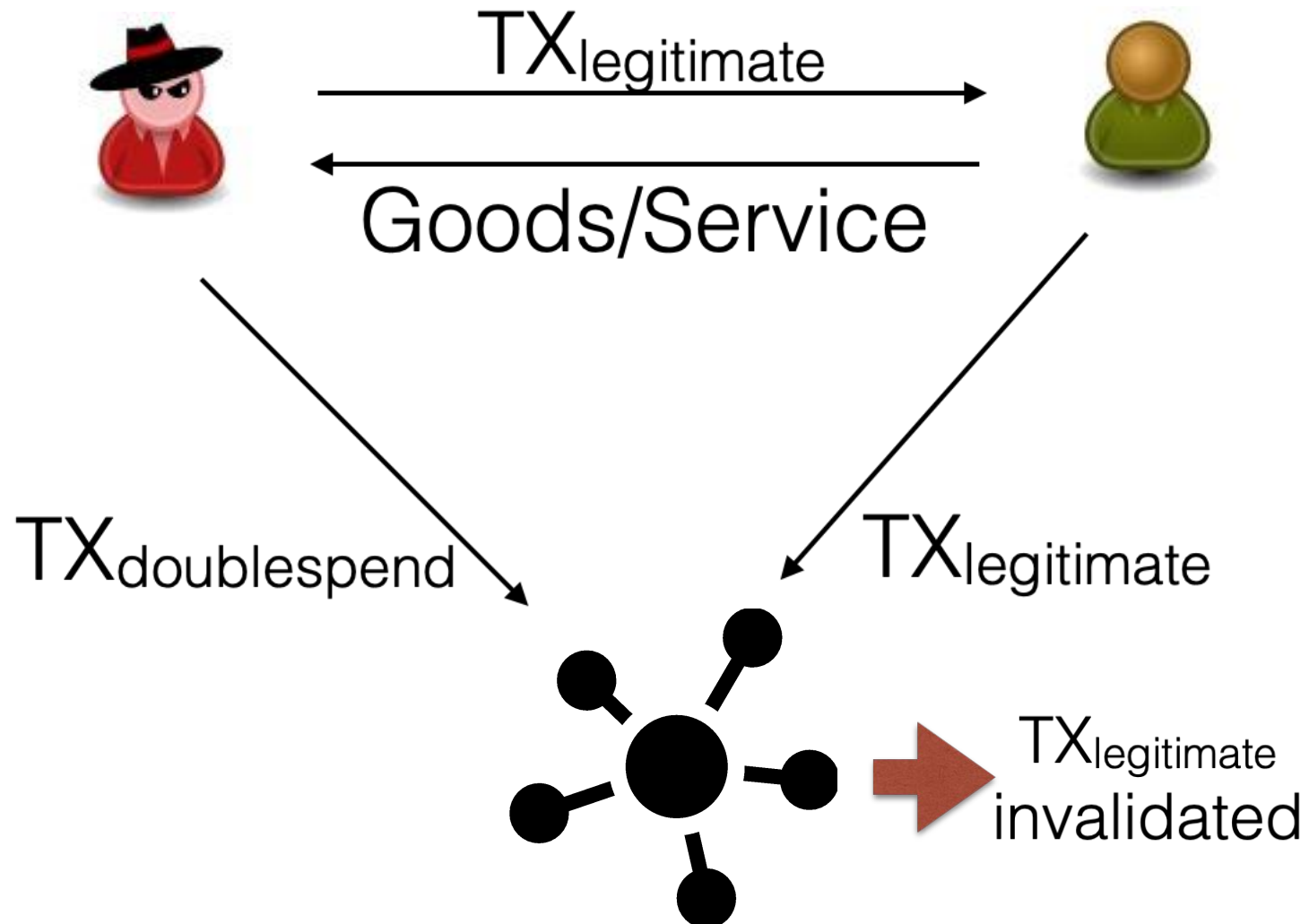
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# Why Blockchain Layer?

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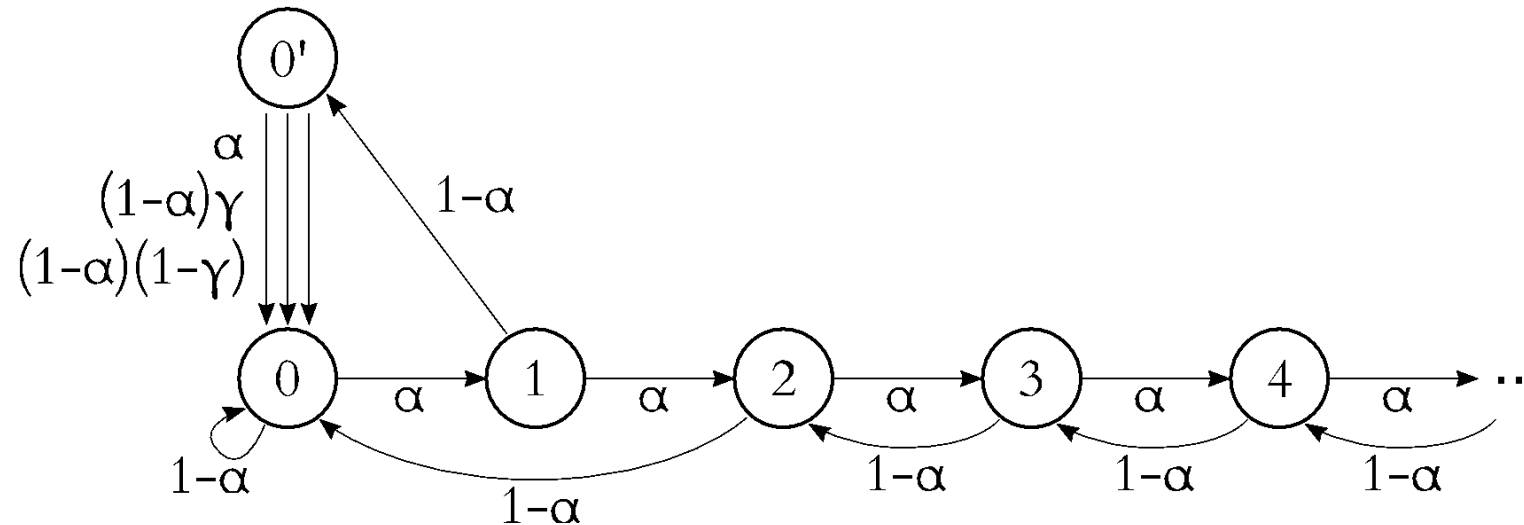
- Double-Spending
- Selfish Mining
- Undercutting
- Bribery

# Double-Spending



# Increasing Mining Advantage with an Eclipse

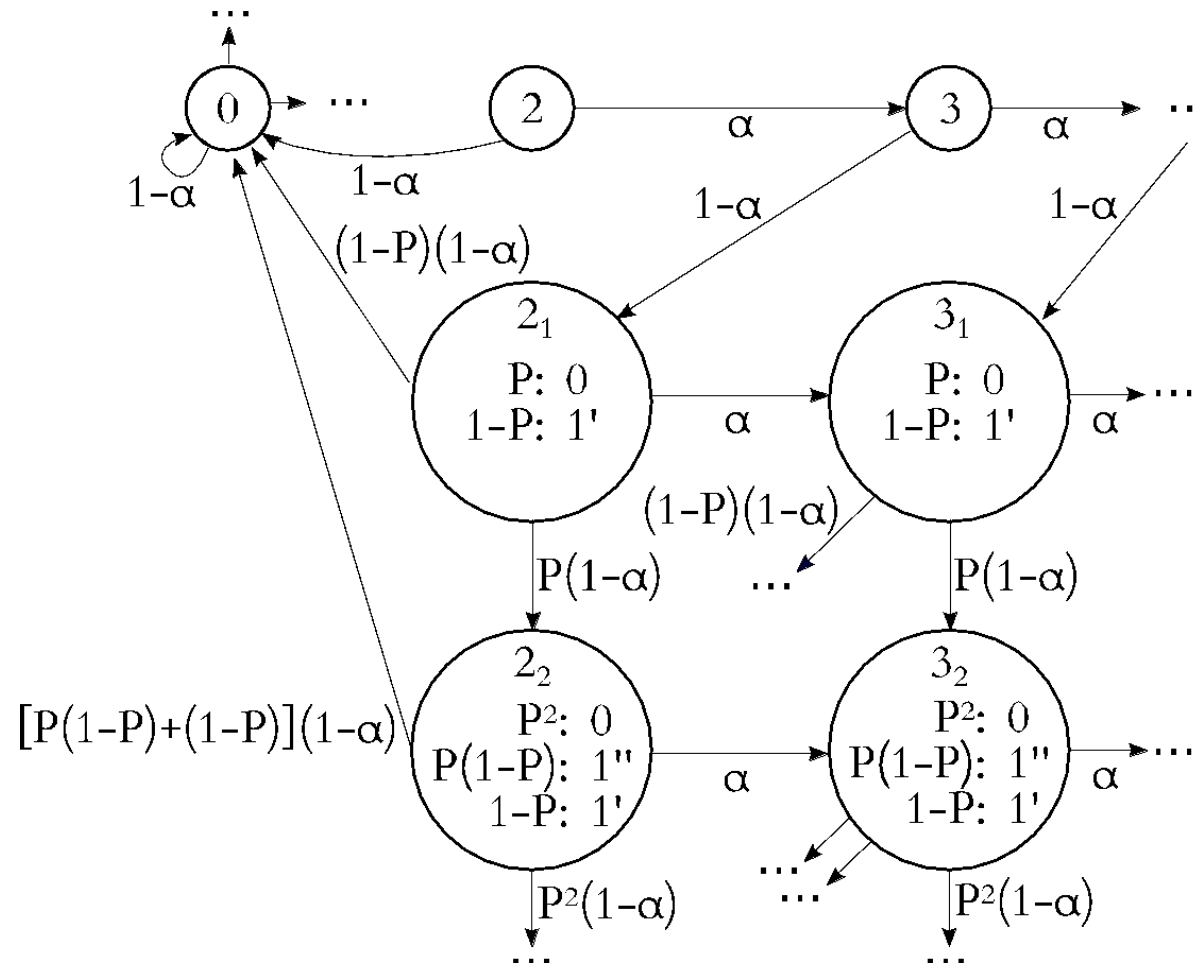
- Idea from Eyal et. al:
  - Instead of publishing, keep a block private
  - Other miners will perform wasteful computations



$\alpha$  : hashing power of adversary

$\gamma$  : propagation parameter

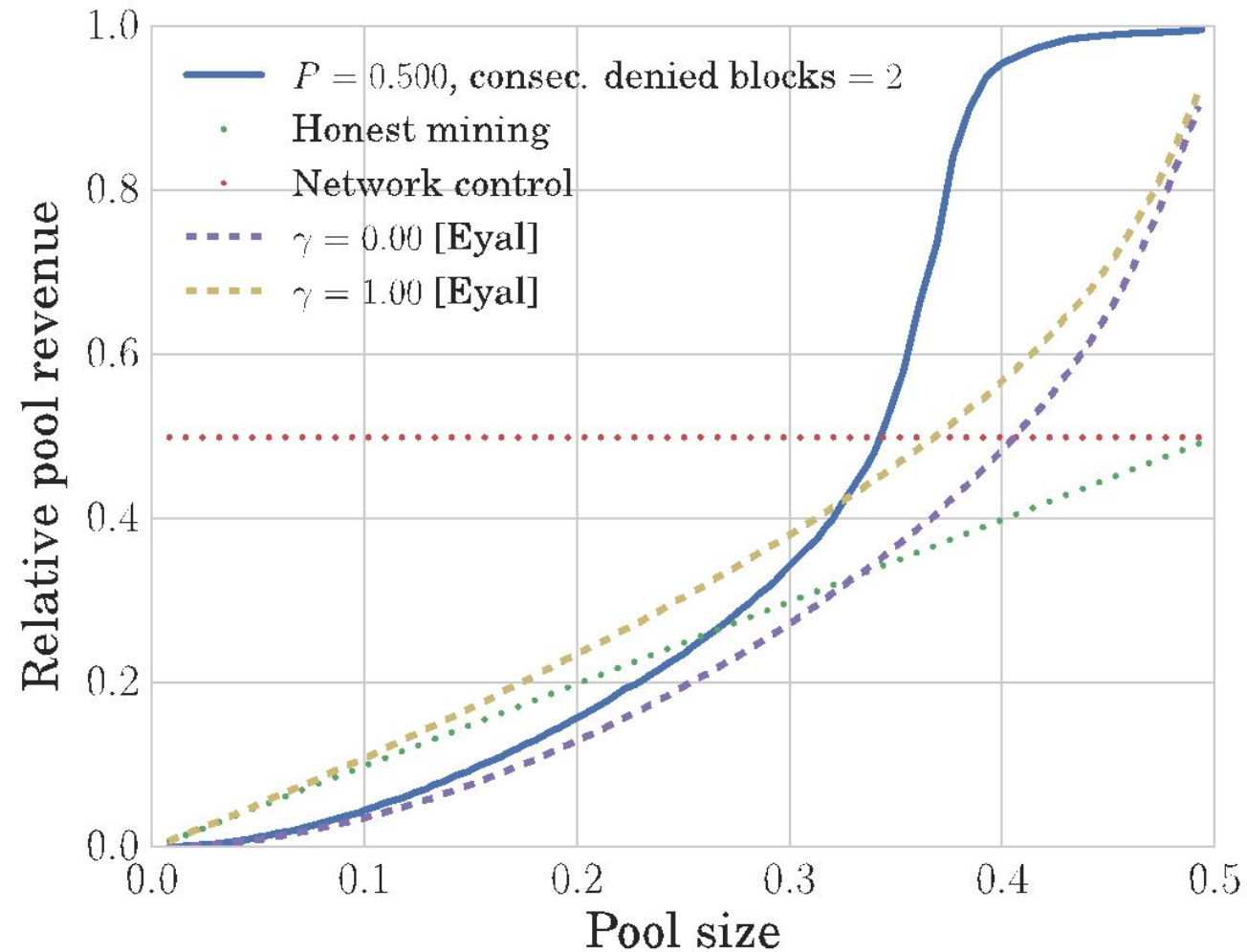
# Increasing Mining Advantage with an Eclipse



P: probability to eclipse a block to a miner



# Increasing Mining Advantage with an Eclipse





# Smart Contract Layer Security

---

# Smart Contract Layer

```
contract Wallet {  
    uint balance = 10;  
  
    function withdraw() {  
        if(balance > 0)  
            msg.sender.call.value(balance) ();  
        balance = 0;  
    }  
}
```

Transfer \$\$\$  
to the caller

- Programs that handle money
  - Executed on a blockchain, written in a high-level language, compiled to VM code
- No patching after release
- What can go wrong?

# The DAO attack

Funds Stolen From the DAO One Year Ago Would be Worth \$1.35bn Today

JP Buntinx June 18, 2017 Crypto. News



33



Etherdice is down for maintenance. We are having troubles with our smart contract and will probably need to invoke the fallback mechanism.

## King of the Ether Throne

An Ethereum ÐApp (a "contract"), living on the blockchain, that will make you a King or Queen, might grant you riches, and will immortalize your name.



### Important Notice

A SERIOUS ISSUE has been identified that can cause monarch compensation payments to not be sent.

DO NOT send payments to the contract previously referenced on this page, or attempt to claim the throne. Refunds will CERTAINLY NOT be made for any payments made after this issue was identified on 2016-02-07.

report

# Security Bug #1: Reentrancy

## User Contract

```
function moveBalance() {  
    wallet.withdraw();  
}  
...  
function () payable {  
    // log payment  
}
```

## Wallet Contract

```
uint balance = 10;  
function withdraw() {  
    if (balance > 0) {  
        logSender.call.value(balance)();  
        balance = 0;  
    }  
}
```

withdraw()

10 ether

Later

withdraw()

no transfer

calls the default  
"payable" function

balance is zeroed  
*after* ether transfer

Can the user contract withdraw more than 10 ether?

# Security Bug #1: Reentrancy

## User Contract

```
function moveBalance() {  
    wallet.withdraw();  
}  
  
function () payable {  
    wallet.withdraw();  
}
```



Calls withdraw()  
before balance  
is set to 0

## Wallet Contract

```
uint balance = 10;  
  
function withdraw() {  
    if (balance > 0)  
        msg.sender.call.value(balance)();  
    balance = 0;  
}
```

balance is zeroed  
*after* ether transfer

An adversary stole 3.6M Ether !

# Security Bug #2: Unprivileged write to storage

## Wallet Contract

```
address owner = ...;

function initWallet(address _owner) {
    owner = _owner;
}

function withdraw(uint amount) {
    if (msg.sender == owner) {
        owner.send(amount);
    }
}
```

Any user may change the wallet's owner

Only owner can send ether

An attacker used a similar bug to *steal \$32M*

# Smart Contract Bug Exercise 1

```
contract Example {  
  
    address public owner;  
    string private mySecret;  
  
    constructor {  
        owner = msg.sender;  
    }  
  
    function setSecret(string _secret) public {  
        require(msg.sender == owner);  
        mySecret = _secret;  
    }  
  
    function getSecret() public returns (string) {  
        require(msg.sender == owner);  
        return mySecret;  
    }  
}
```

Any variable is readable on the public Ethereum blockchain. Declaring a variable private only restricts the automatic creation of getter for that variable, but does not hide it.

Hint: who would be able to read mySecret?



# Smart Contract Bug Exercise 2

```
contract Vulnerable {  
  
    mapping(address => bool) authorized;  
    mapping(address => uint) balances;  
  
    function refund(uint amount) public {  
        require(authorized[msg.sender]);  
        require(amount <= balances[msg.sender]);  
  
        msg.sender.call.value(amount)("");  
        balances[msg.sender] -= amount;  
    }  
}
```

The code is vulnerable to a **reentrancy attack**.  
The balance of the *msg.sender* is only updated after a transfer is made. If the *msg.sender* is a contract and has a fallback function that calls into the contract again, the *msg.sender* can deplete the contract of the funds.

Hint: who can be *msg.sender*?

# Smart Contract Bug Exercise 2

```
contract Vulnerable {
    ... // vulnerable as the previous example
}

contract Exploit {

    Vulnerable v;

    function register(address contract) public {
        v = Vulnerable(contract);
    }

    function exploit() public {
        // your code here
    }

    // your code here
}
```

Hint: check the previous example

# Smart Contract Bug Exercise 2 - Solution

```
contract Vulnerable {
    ... // vulnerable as the previous example
}

contract Exploit {

    Vulnerable v;

    function register(address contract) public {
        v = Vulnerable(contract);
    }

    function exploit() public {
        v.refund(1);
    }

    function () public {
        v.refund(1);
    }
}
```

# More smart contract security bugs

---



Unexpected ether flows



Insecure coding, such as unprivileged writes (e.g., *Multisig Parity bug*)



Use of unsafe inputs (e.g., reflection, hashing, ...)



Reentrant method calls (e.g., *DAO bug*)

# More smart contract security bugs

Known Attacks

The following is a list of known attacks which you should be aware of, and defend against when writing smart contracts.

## Reentrancy

One of the major dangers of [calling external contracts](#) is that they can take over the control flow, and make changes to your data that the calling function wasn't expecting. This class of bug can take many forms, and both of the major bugs that led to the DAO's collapse were bugs of this sort.

### Reentrancy on a Single Function

The first version of this bug to be noticed involved functions that could be called repeatedly, before the first invocation of the function was finished. This may cause the different invocations of the function to interact in destructive ways.

```
// INSECURE
mapping (address => uint) private userBalances;

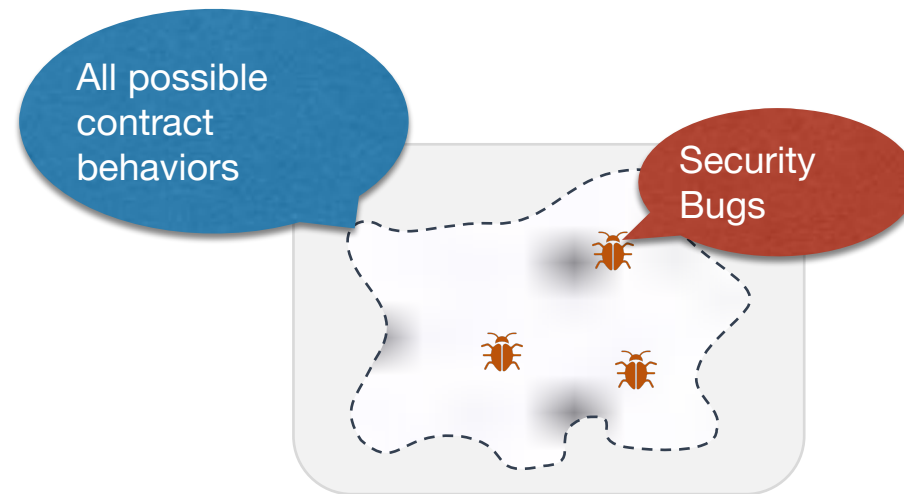
function withdrawBalance() public {
    uint amountToWithdraw = userBalances[msg.sender];
    (bool success, ) = msg.sender.call.value(amountToWithdraw)(""); // At this point
    require(success);
    userBalances[msg.sender] = 0;
}
```

Since the user's balance is not set to 0 until the very end of the function, the second (and later)

### Table of contents

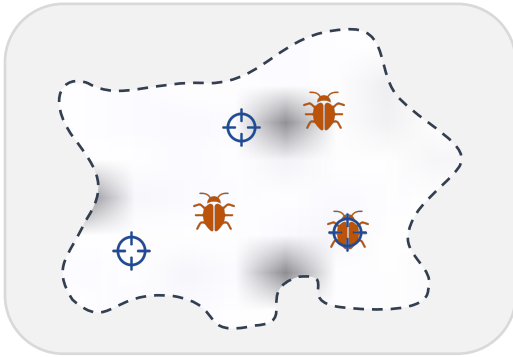
- Reentrancy
  - Reentrancy on a Single Function
  - Cross-function Reentrancy
  - Pitfalls in Reentrancy Solutions
- Front-Running
  - Taxonomy
    - Displacement
    - Insertion
    - Suppression
  - Mitigations
- Timestamp Dependence
- Integer Overflow and Underflow
- DoS with (Unexpected) revert
- DoS with Block Gas Limit
  - Gas Limit DoS on a Contract via Unbounded Operations
  - Gas Limit DoS on the Network via Block Stuffing
- Insufficient gas griefing
- Forcibly Sending Ether to a Contract
- Deprecated/historical attacks
  - Call Depth Attack (deprecated)
  - Constantinople Reentrancy Attack
- Other Vulnerabilities

# Automated security analysis



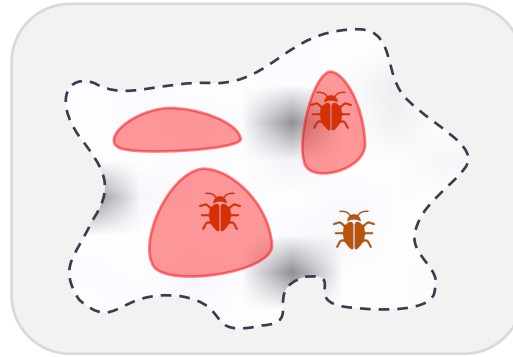
Problem: Cannot enumerate all possible contract behaviors...

# Automated security analysis – Existing solutions



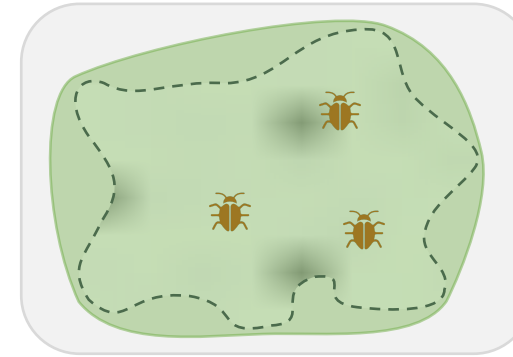
Testing

Easy to implement, but  
very limited guarantees



Dynamic analysis  
Symbolic execution

Better than testing, but  
can still miss vulnerabilities



Static analysis  
Formal verification

Strong guarantees, but many  
false positives



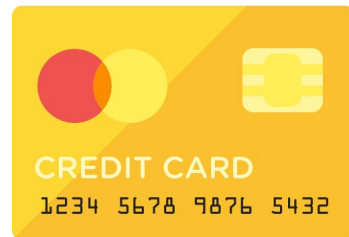
# DeFi Flash Loan „Attacks“

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# Flash Loan Attacks

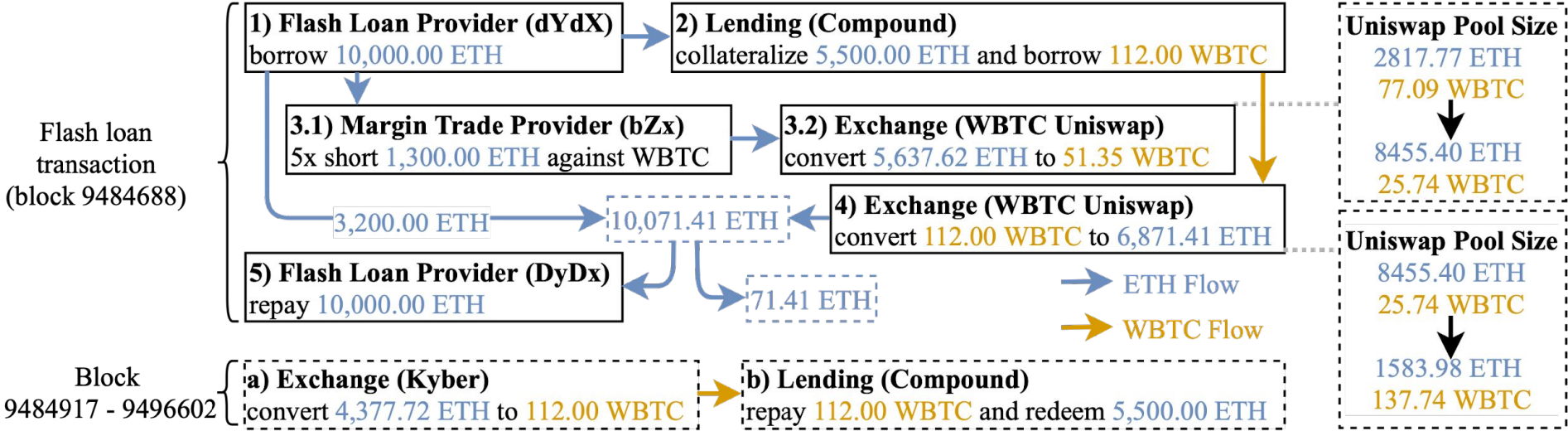
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+



# bZx - Pump and Arbitrage Attack – February 2020



**Input: 130 USD gas**

**Output: 350,000 USD**

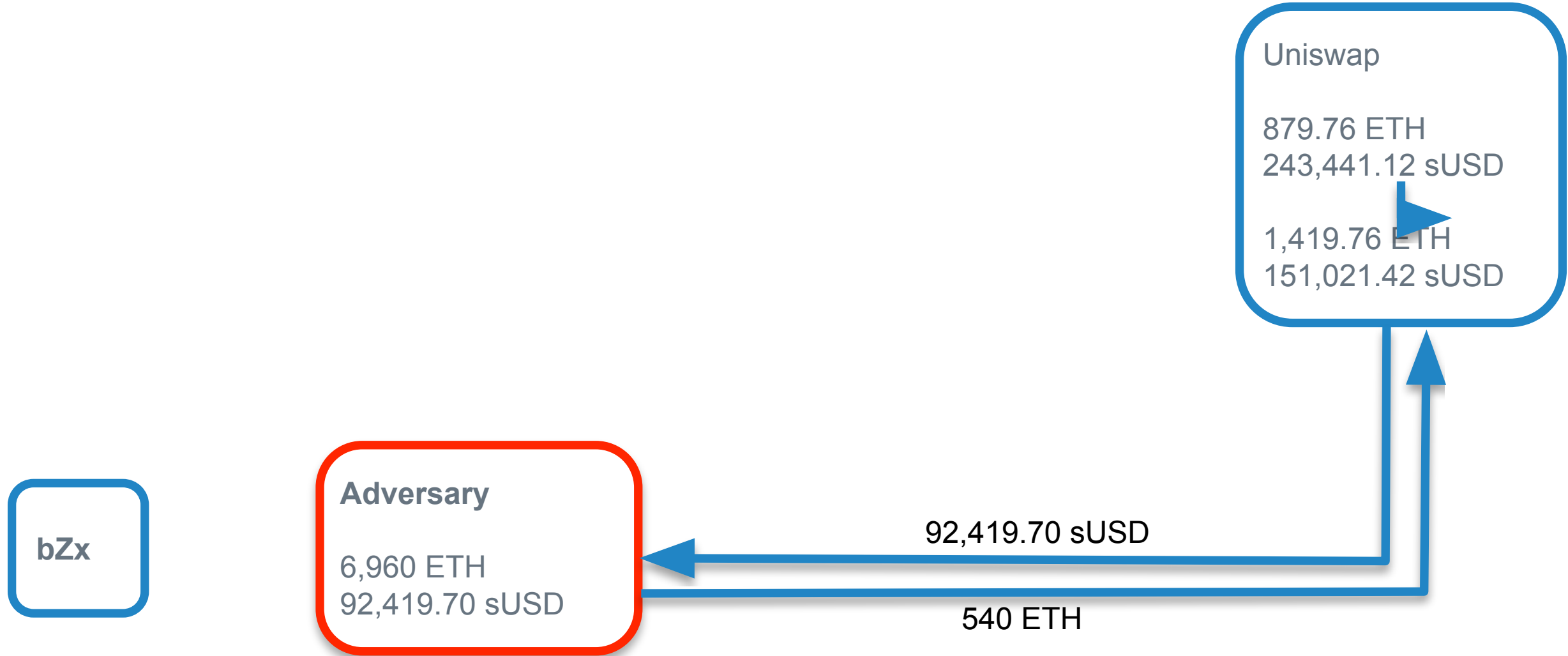
**Optimal: 830,000 USD**

# bZx – Oracle manipulation – February 2020

---

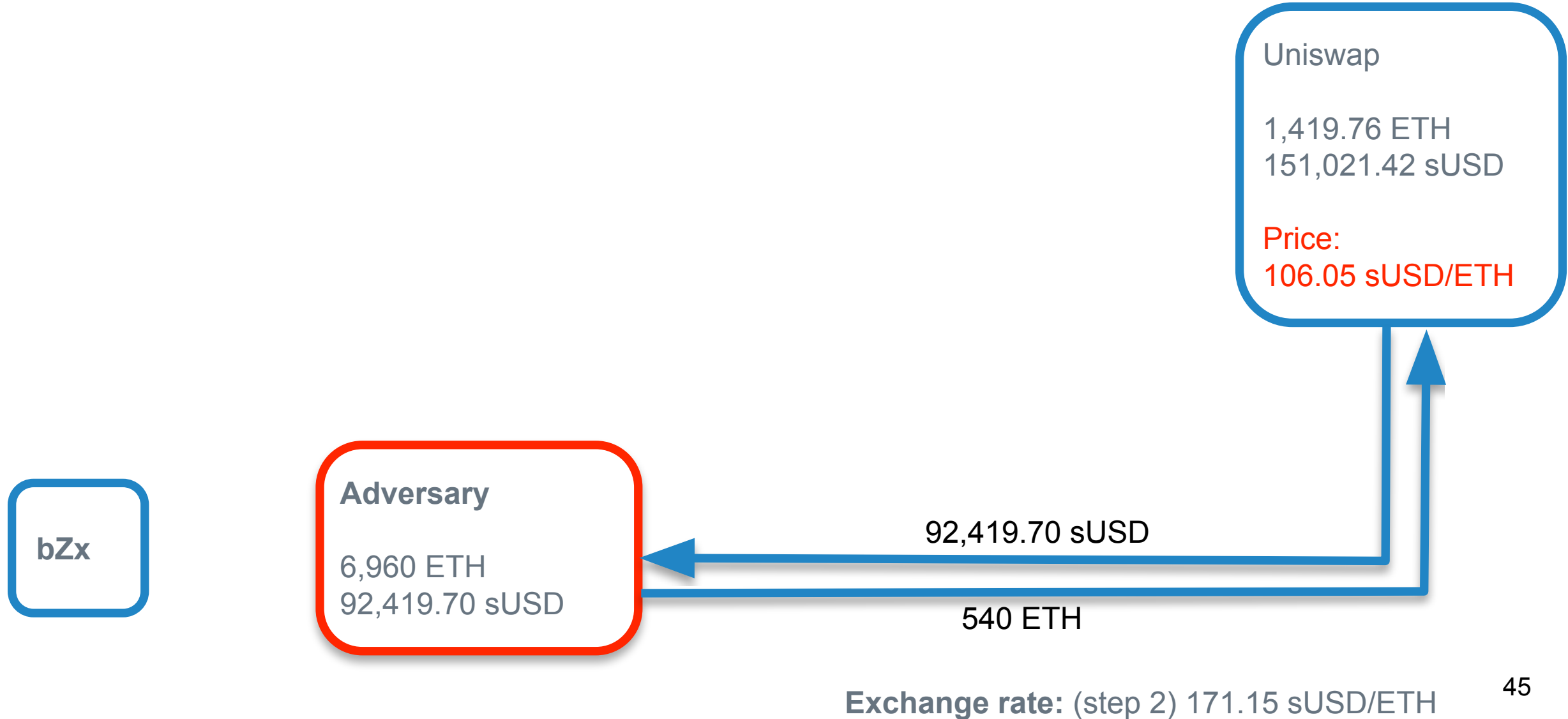


# bZx – Oracle manipulation – February 2020

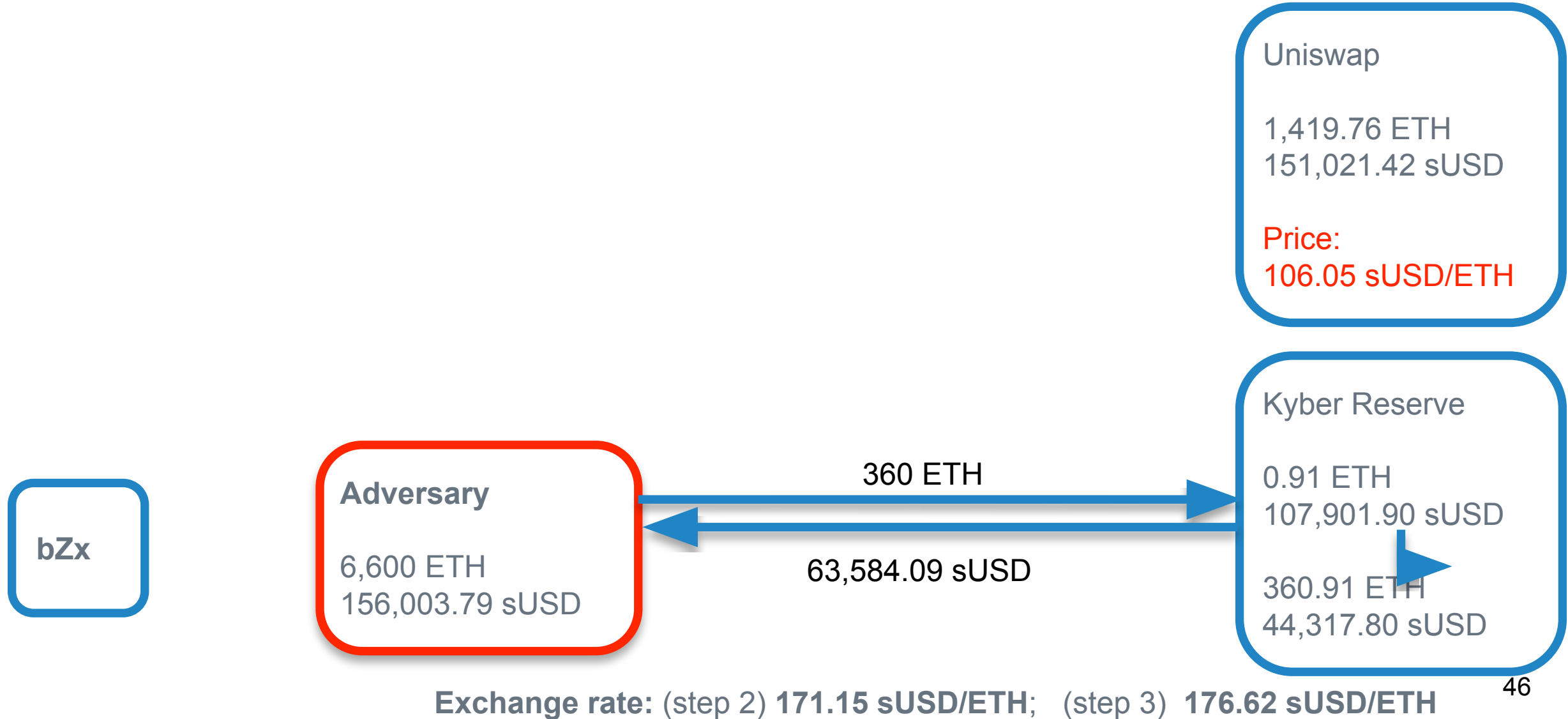


Exchange rate: (step 2) 171.15 sUSD/ETH

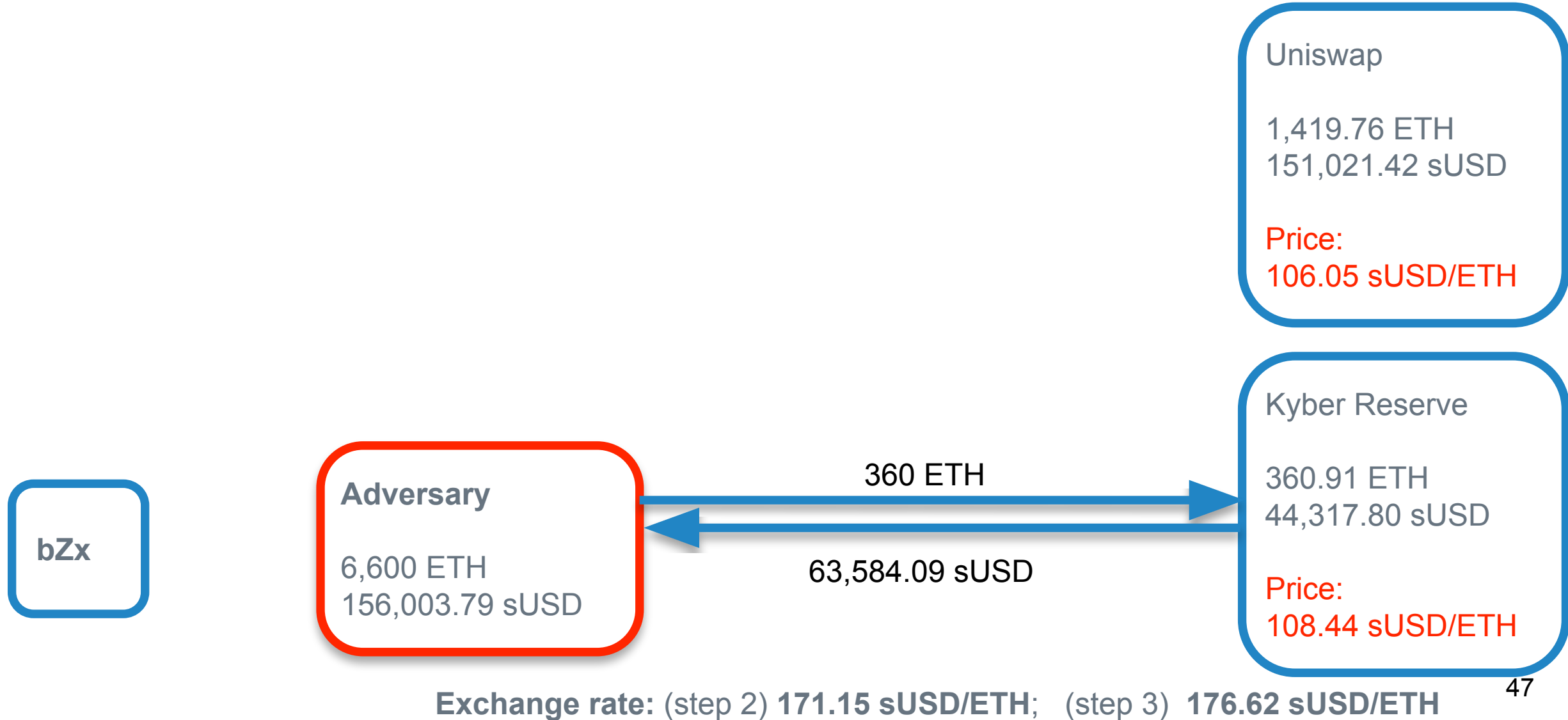
# bZx – Oracle manipulation – February 2020



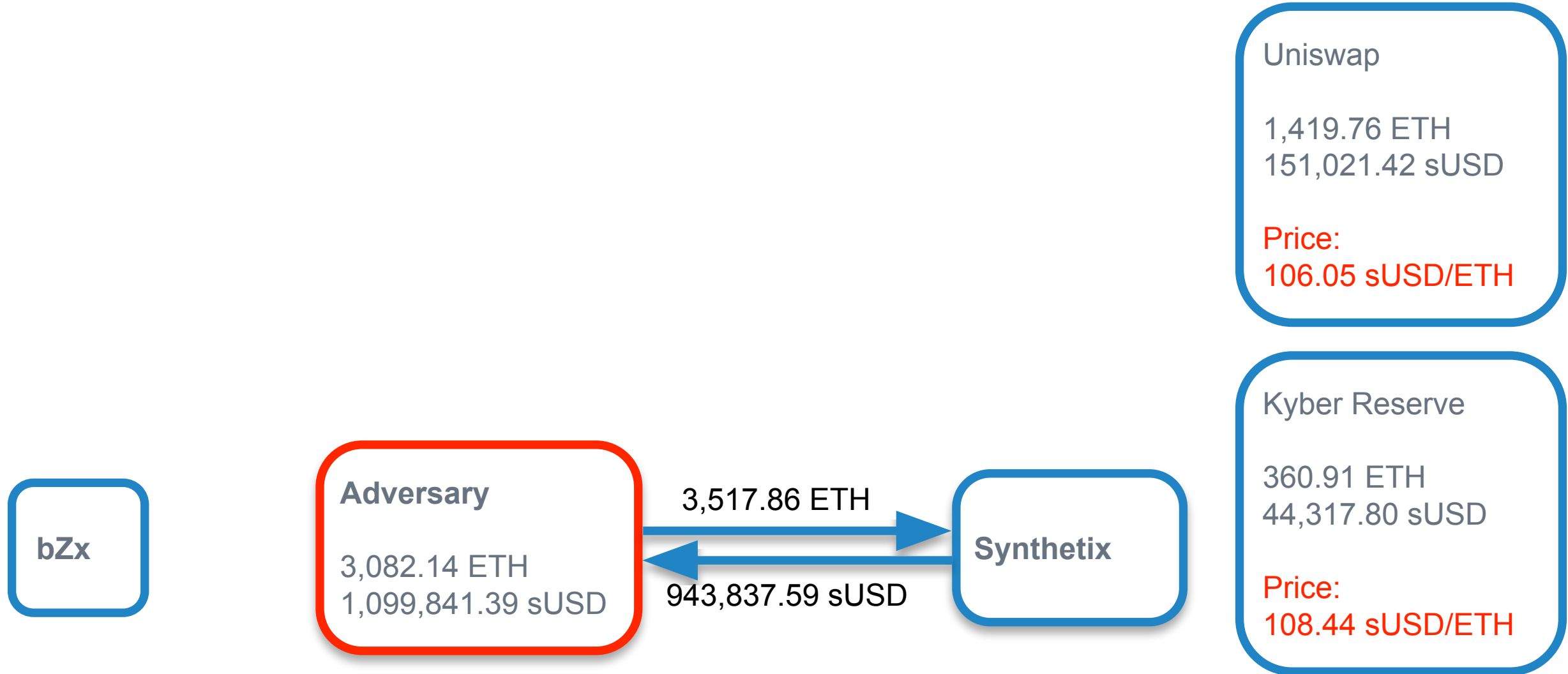
# bZx – Oracle manipulation – February 2020



# bZx – Oracle manipulation – February 2020



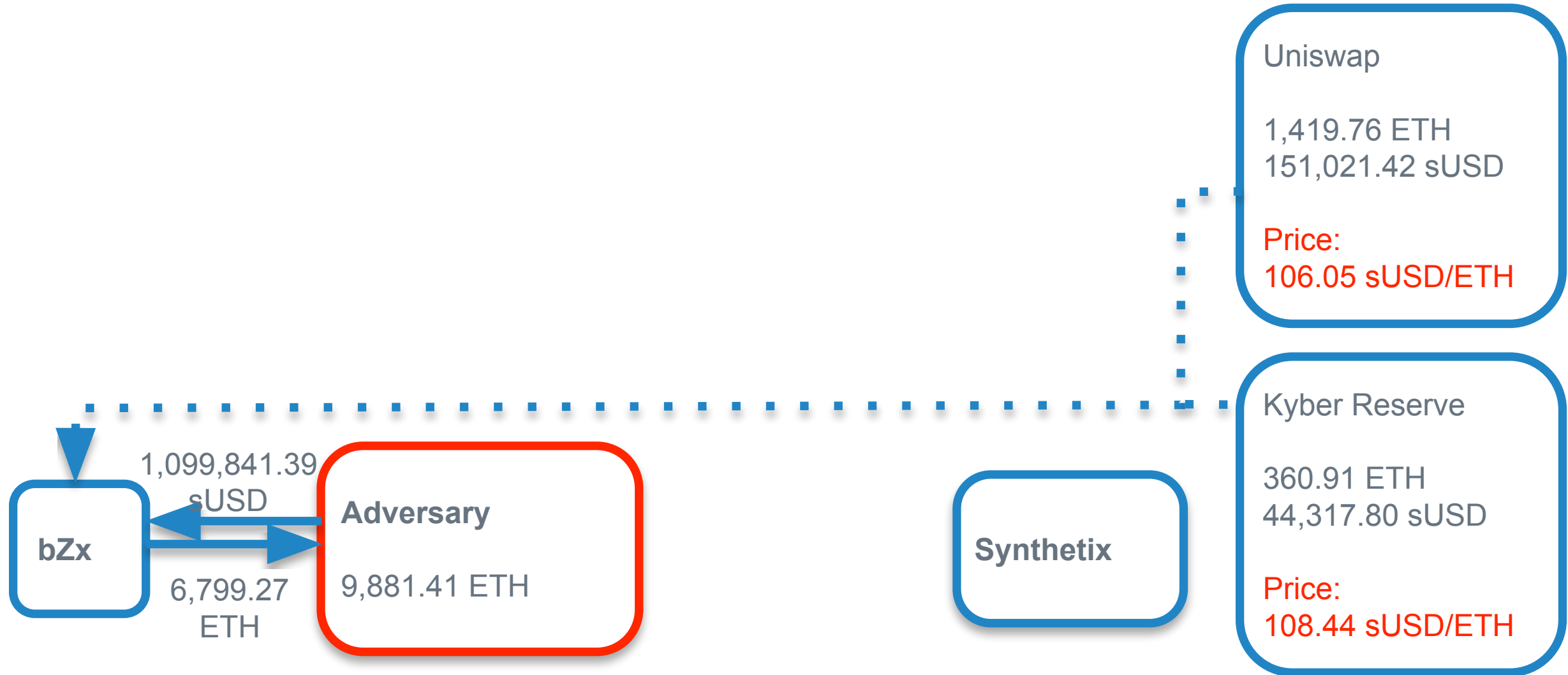
# bZx – Oracle manipulation – February 2020



Exchange rate: (step 2) 171.15 sUSD/ETH; (step 3) 176.62 sUSD/ETH; (step 4) 268.30 sUSD/ETH

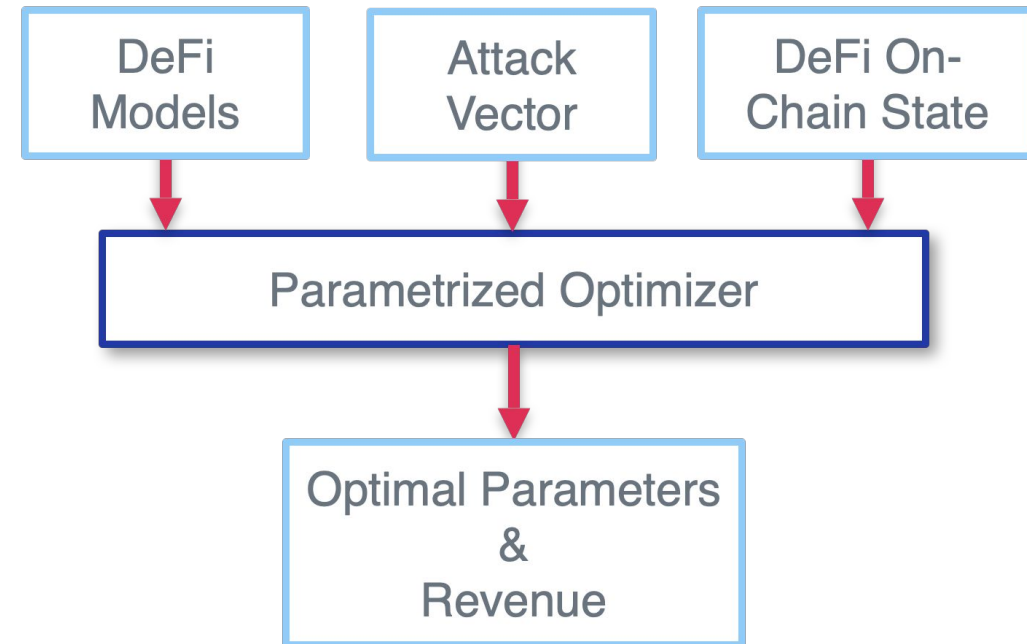


# bZx – Oracle manipulation – February 2020



# Constrained Optimization Framework

- Formulate DeFi actions in models
  - Constant product AMM:  $\Delta y = y - \frac{xy}{x+\Delta x}$
- Construct a constrained optimization problem based on the attack vector
  - Objective function: outcome profit
- Fetch the on-chain state that the



# Optimizing the bZx attack 2

- Borrow  $X$  ETH (bZx flash loan)
  - Convert  $p1$  ETH to  $f1(p1)$  sUSD (Uniswap)
  - Convert  $p2$  ETH to  $f2(p2)$  sUSD (Kyber)
  - Deposit  $p3$  ETH for  $f3(p3)$  sUSD (Synthetix)
  - Collateralize  $z$  sUSD to borrow  $g(z)$  ETH
    - $z=f1(p1)+f2(p2)+f3(p3)$
- Repay  $X$  ETH (bZx flash loan)
- Objective:  $o=g(f1(p1)+f2(p2)+f3(p3))-X$ 
  - s.t.  $p1+p2+p3<X$

# Optimizing the bZx attack 2

---

- Sequential Least Squares Programming (SLSQP)
  - SciPy
- Ubuntu 18.04.2, 16 CPU cores, 32 GB RAM
- Validation by concrete execution
  - Execution on the real blockchain state



# Sandwich Attacks

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# AMM – Automated Market Maker

---

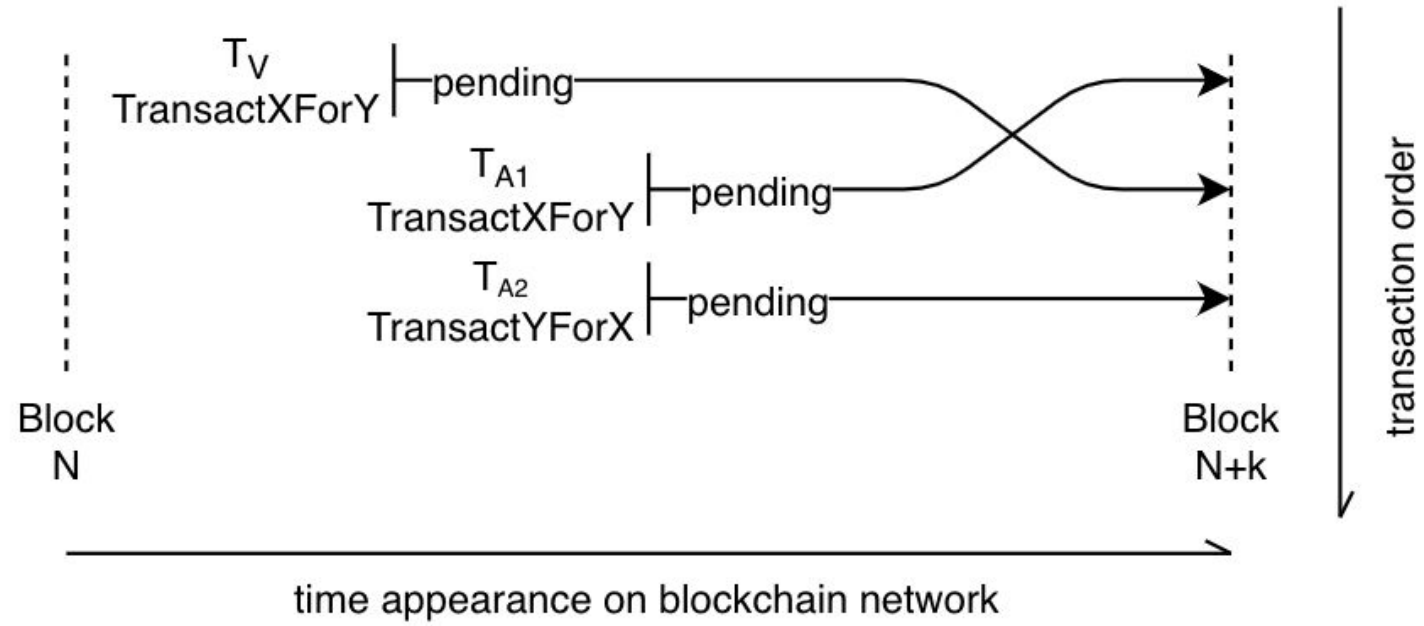
$$x \times y = k$$

Asset X  
quantity

Asset Y  
quantity

constant

# Sandwich Attack



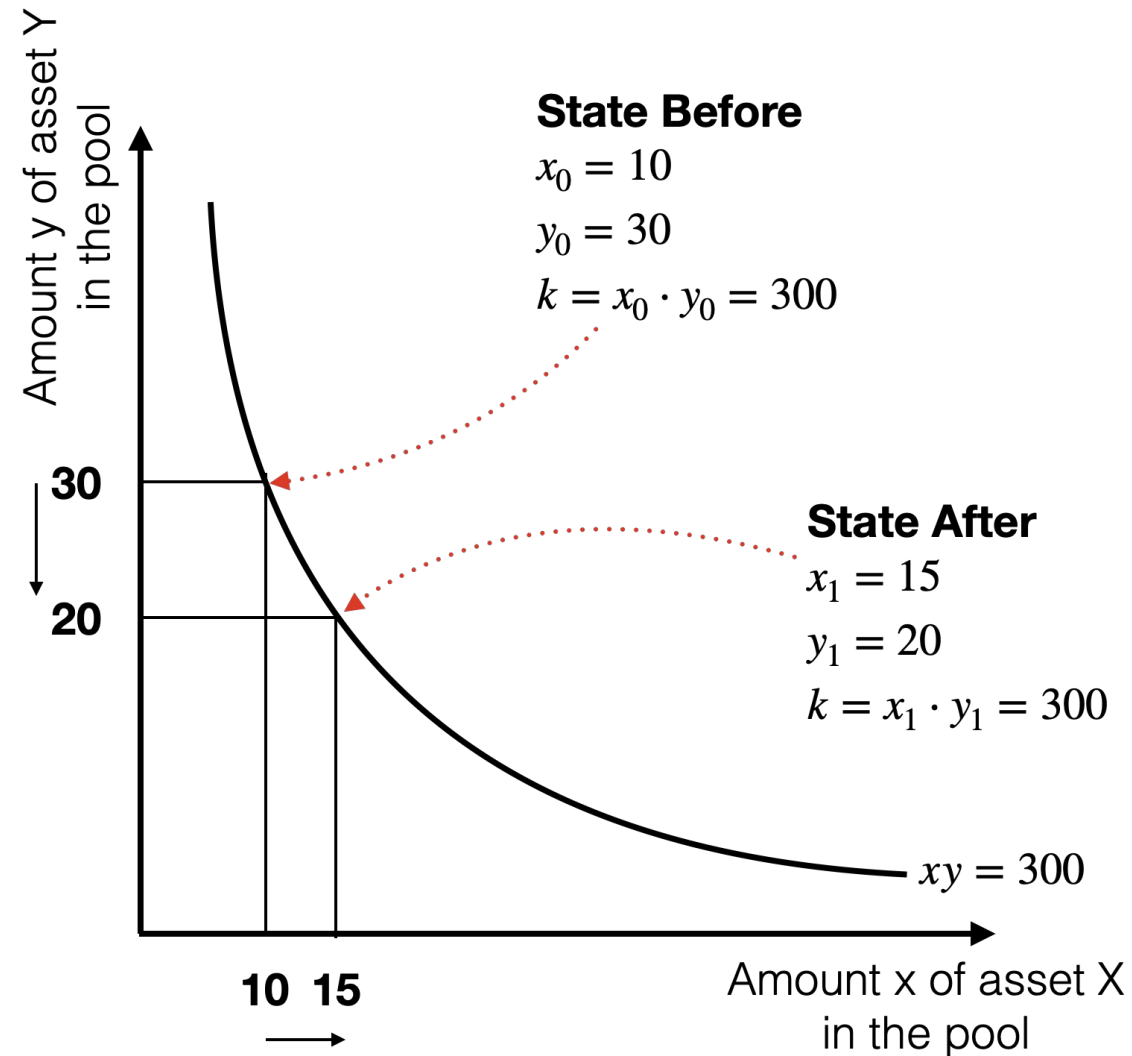
# AMM – Constant product formula

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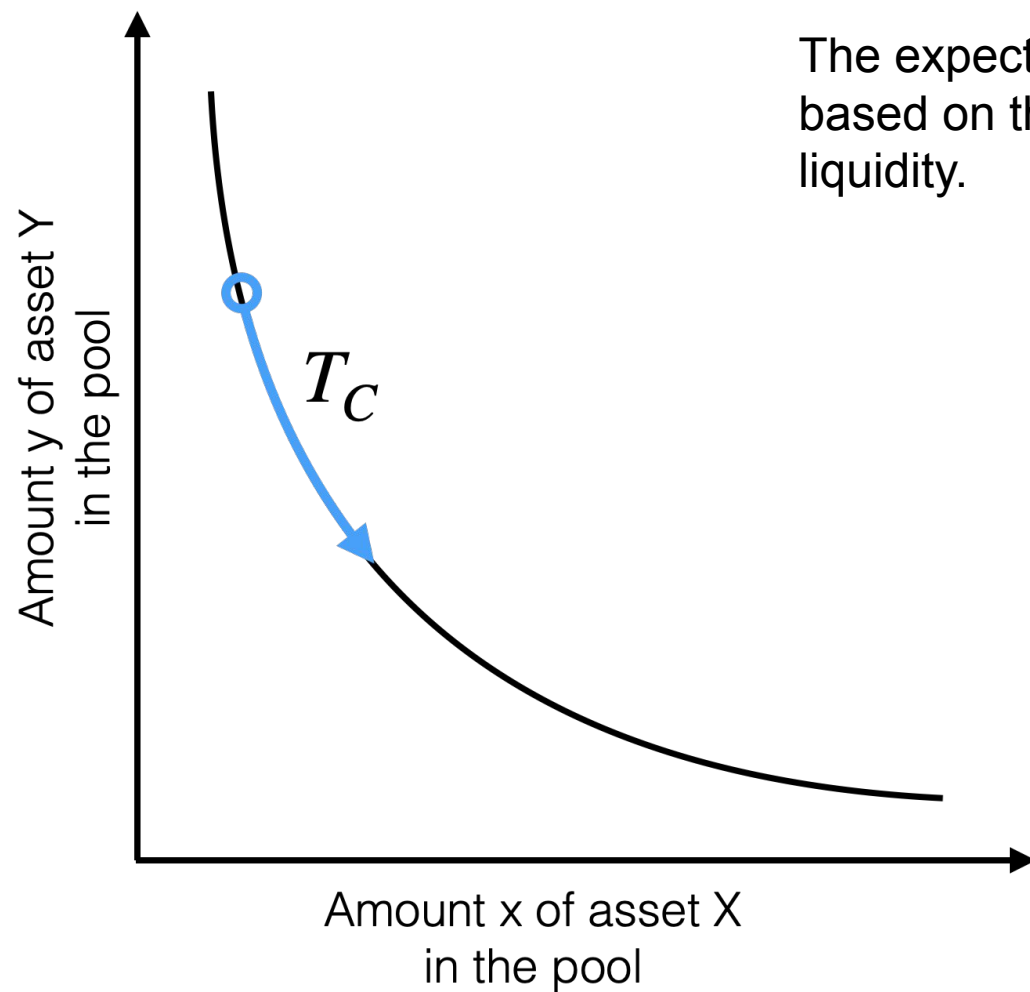




# AMM – Constant product formula

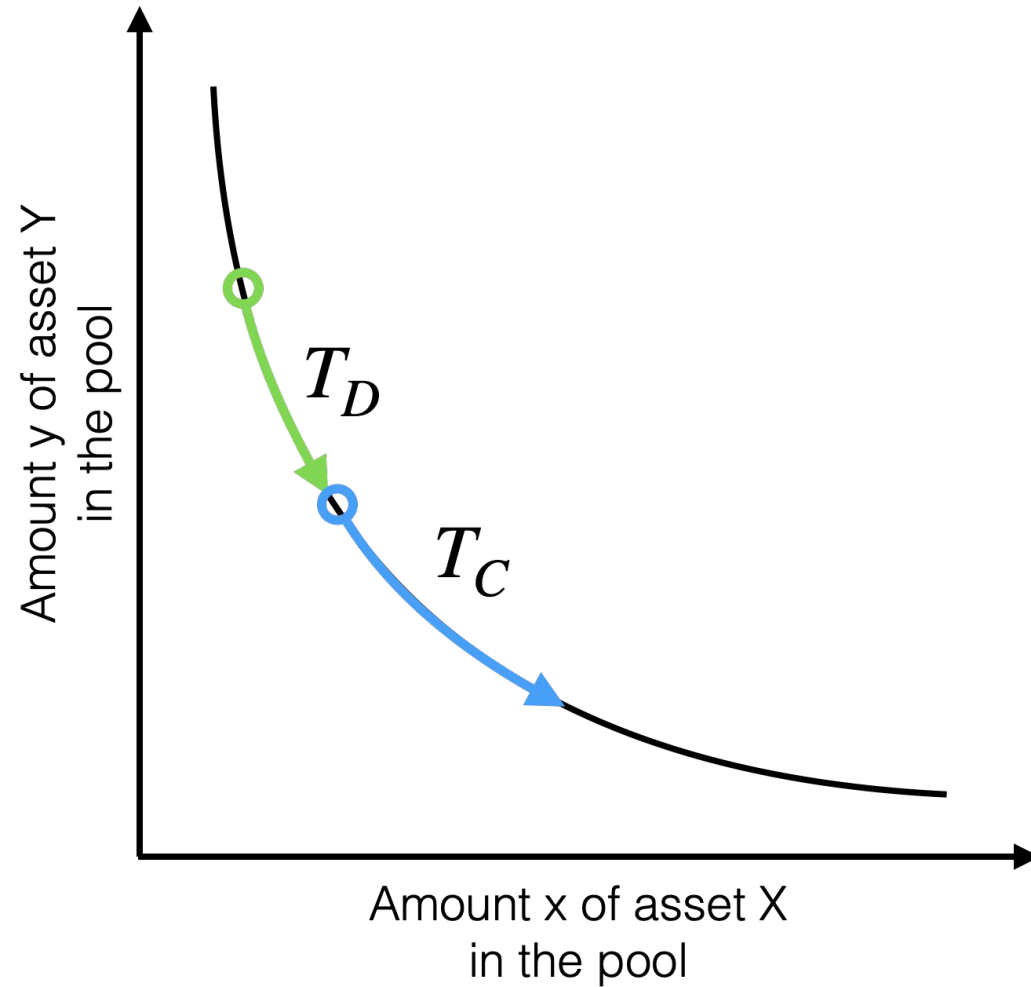


# Expected Slippage

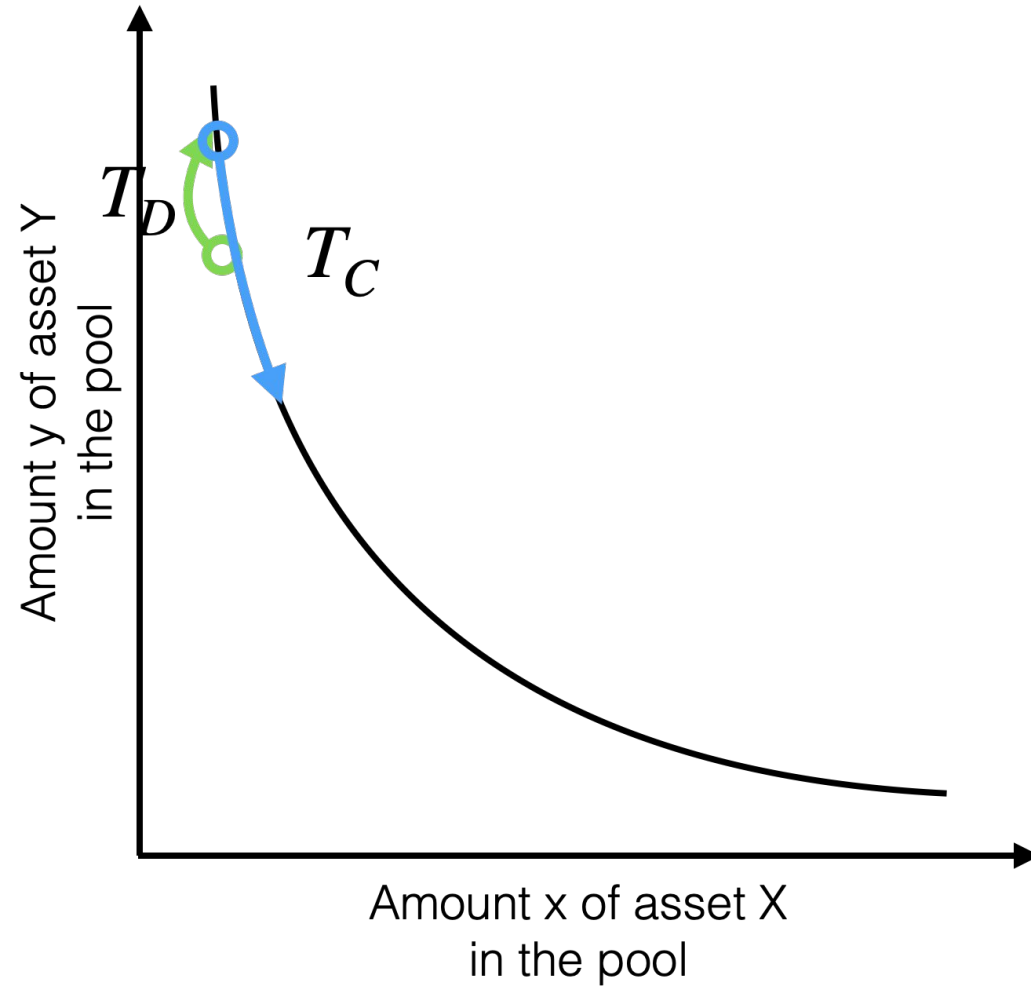


The expected increase or decrease in price based on the trading volume and available liquidity.

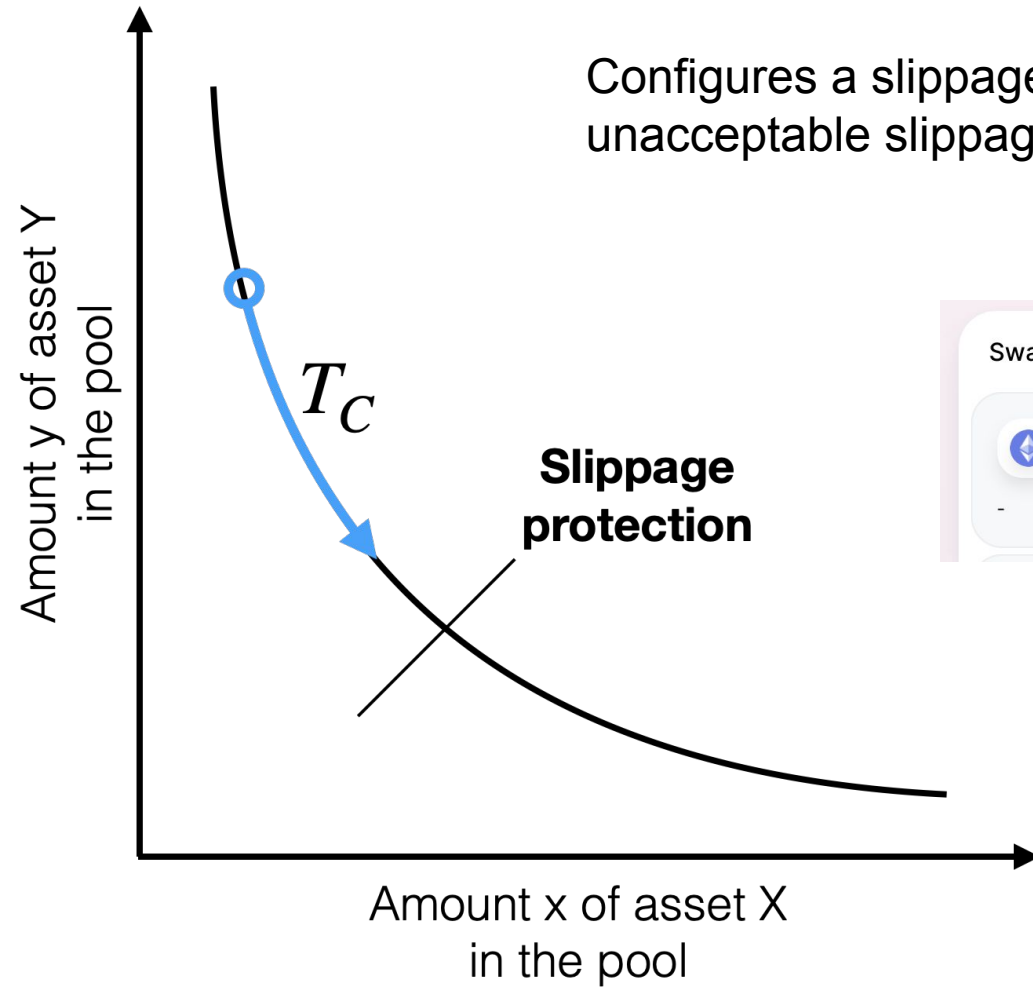
# Unexpected Slippage -> Worse Execution Price



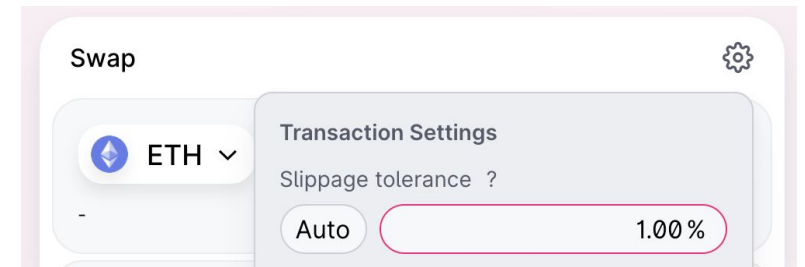
# Unexpected Slippage -> Better Execution Price



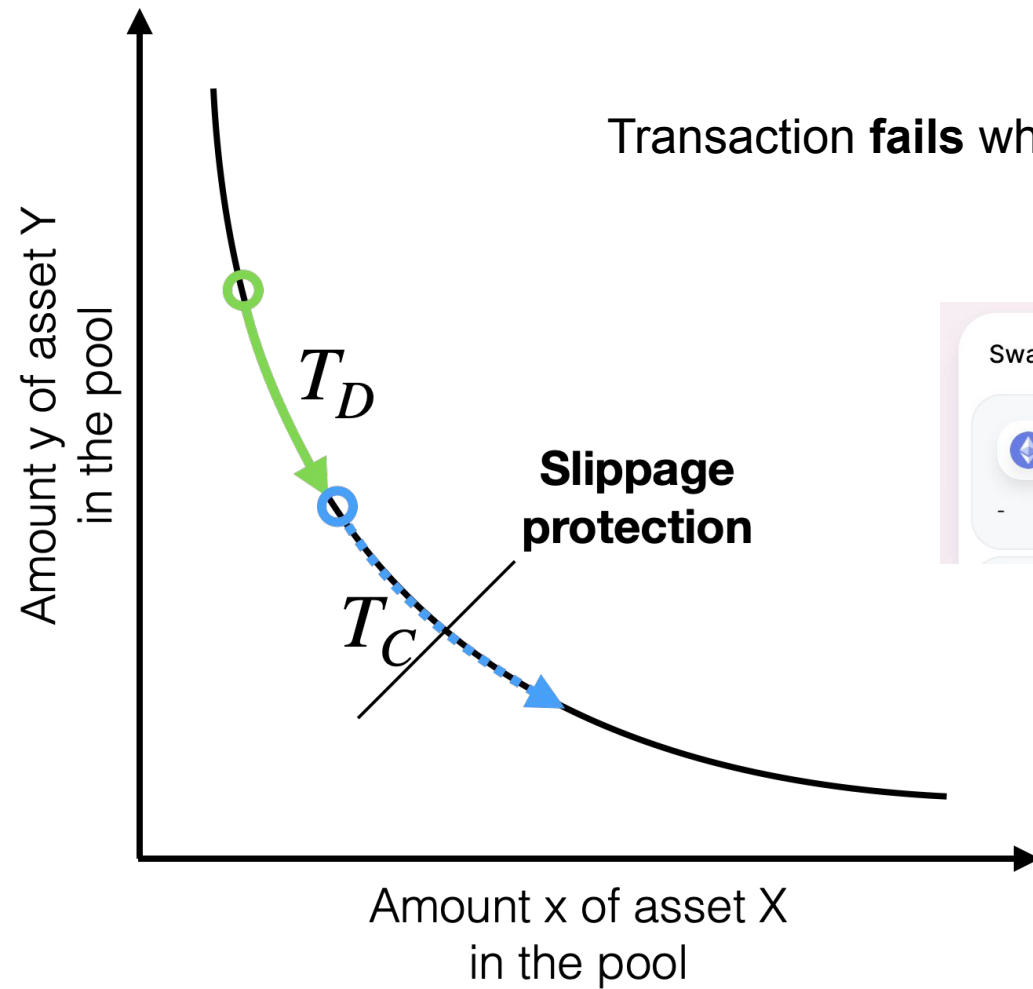
# Slippage Protection



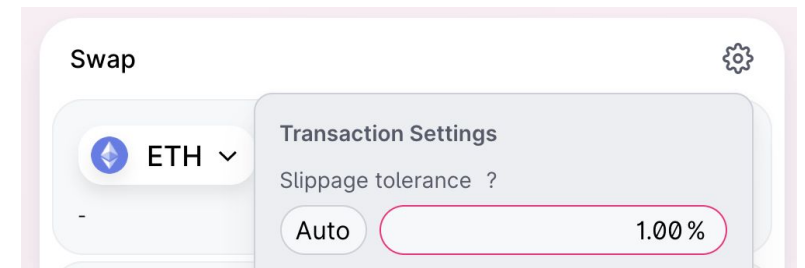
Configures a slippage protection threshold to prevent unacceptable slippage



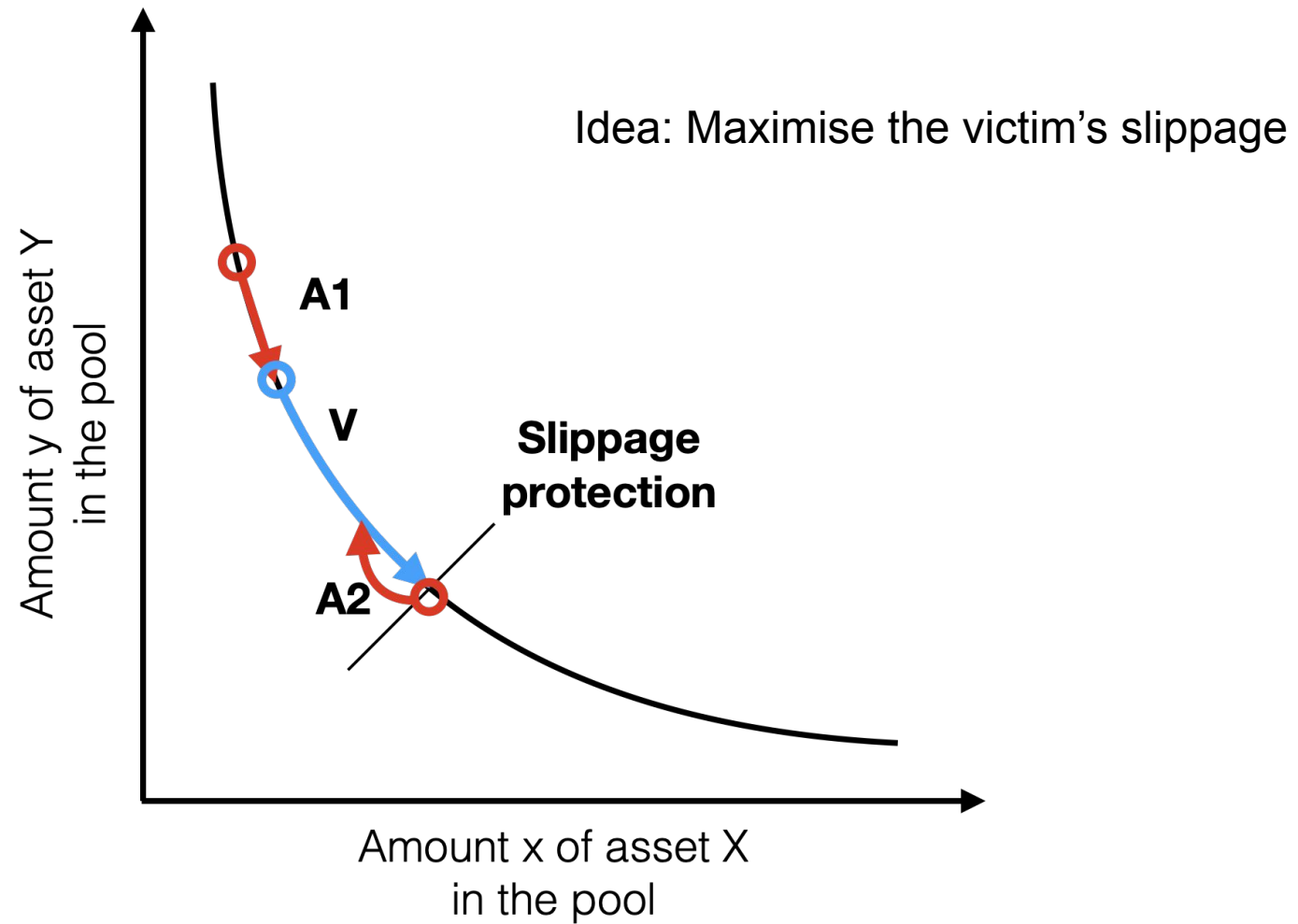
# Slippage Protection



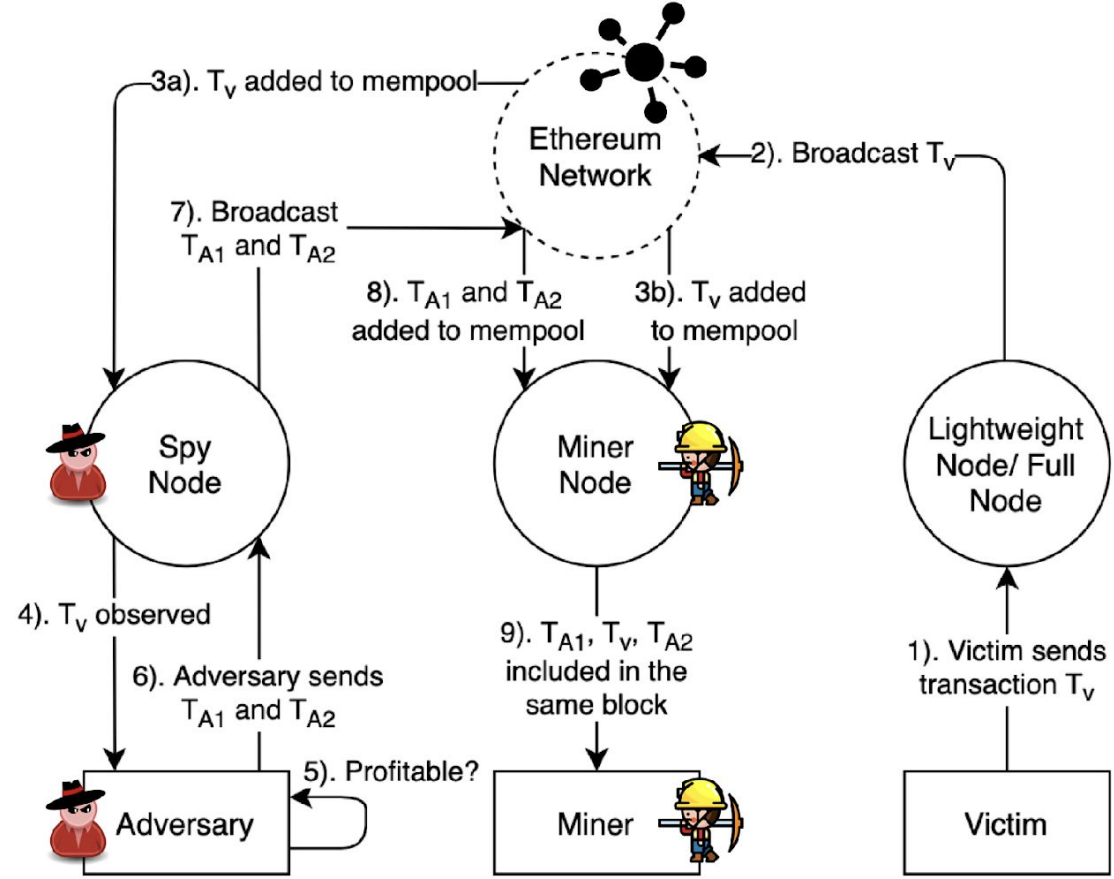
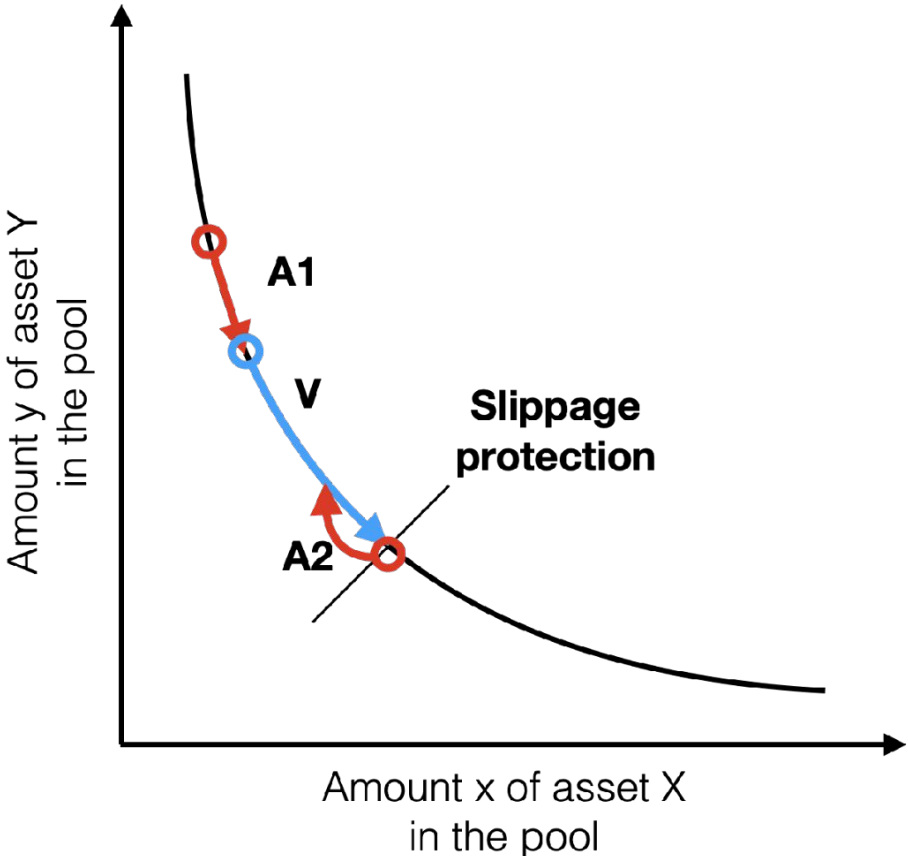
Transaction **fails** when crossing the slippage limit.



# Sandwich Attack Against Taker

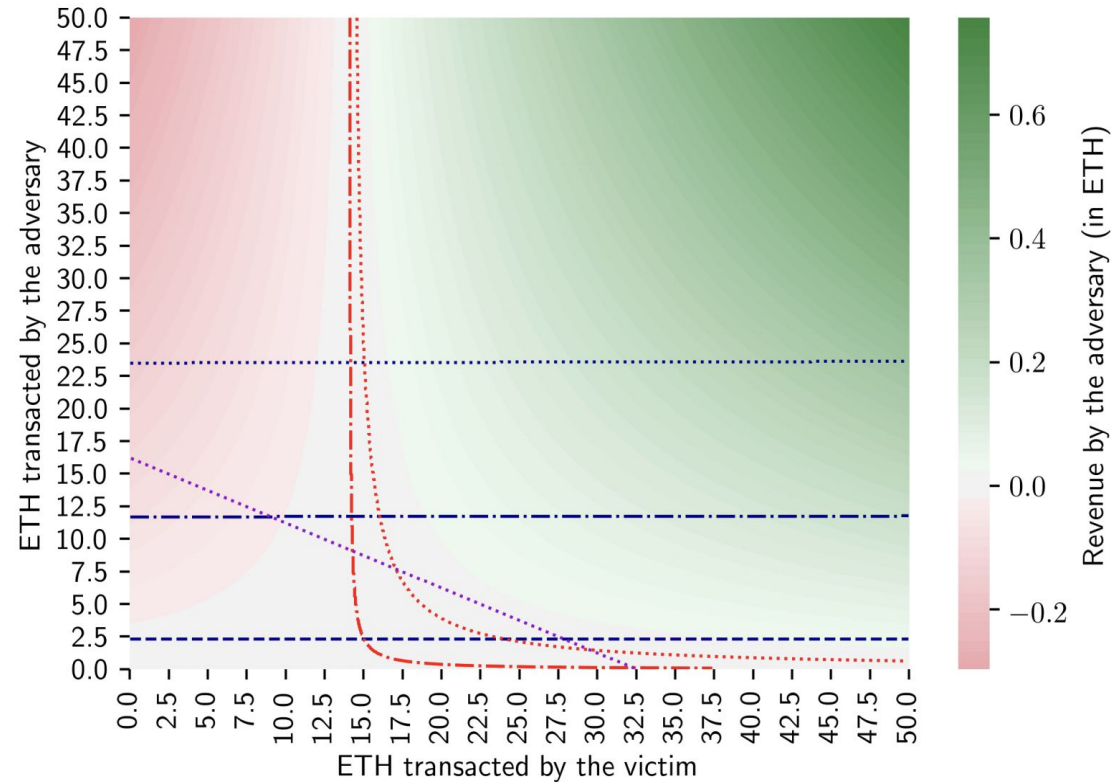


# Network layer + DeFi protocol layer





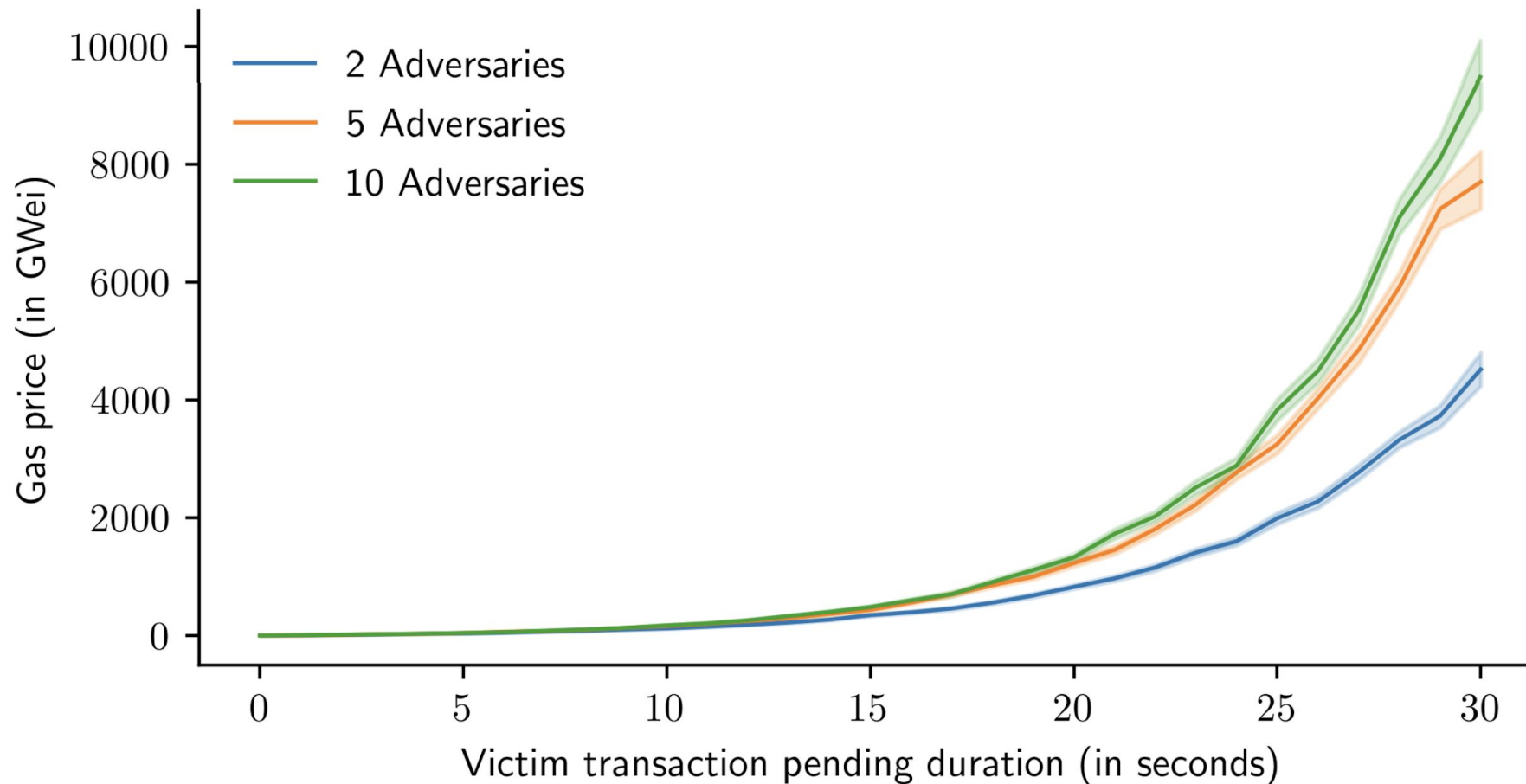
# Sandwich attack profitability



- ..... adversary breaks even at 0.01 ETH transaction fee
- .-.- adversary breaks even at 0.001 ETH transaction fee
- ..... 1% total(expected + unexpected) slippage for the victim
- ..... 1% unexpected slippage for the victim
- .-.- 0.5% unexpected slippage for the victim
- ..... 0.1% unexpected slippage for the victim

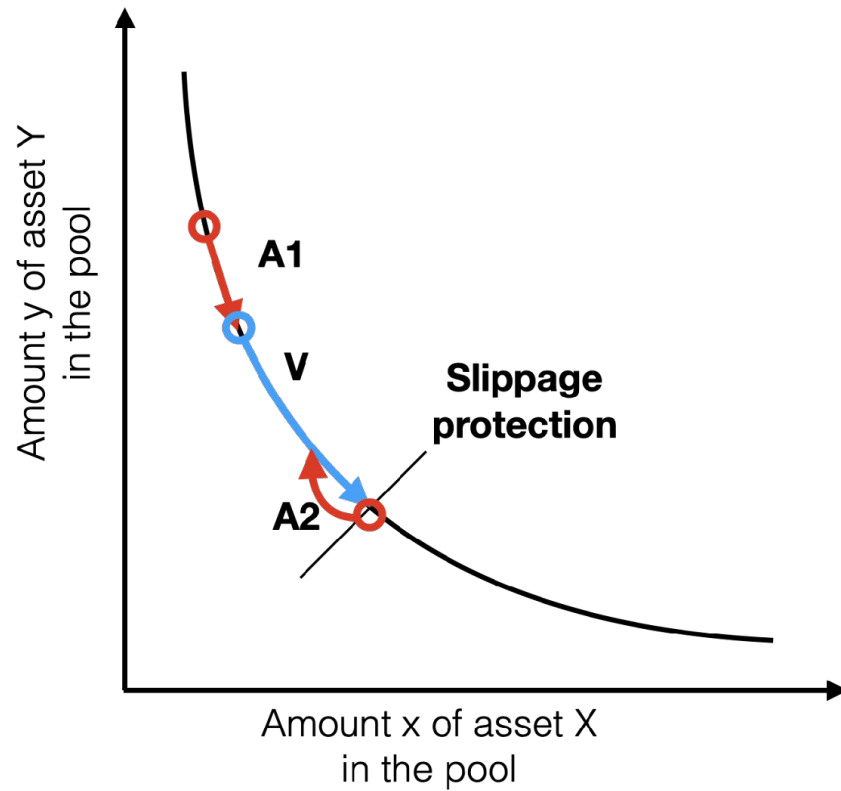
# Multiple Adversaries

Break-even of the attacker becomes harder to attain

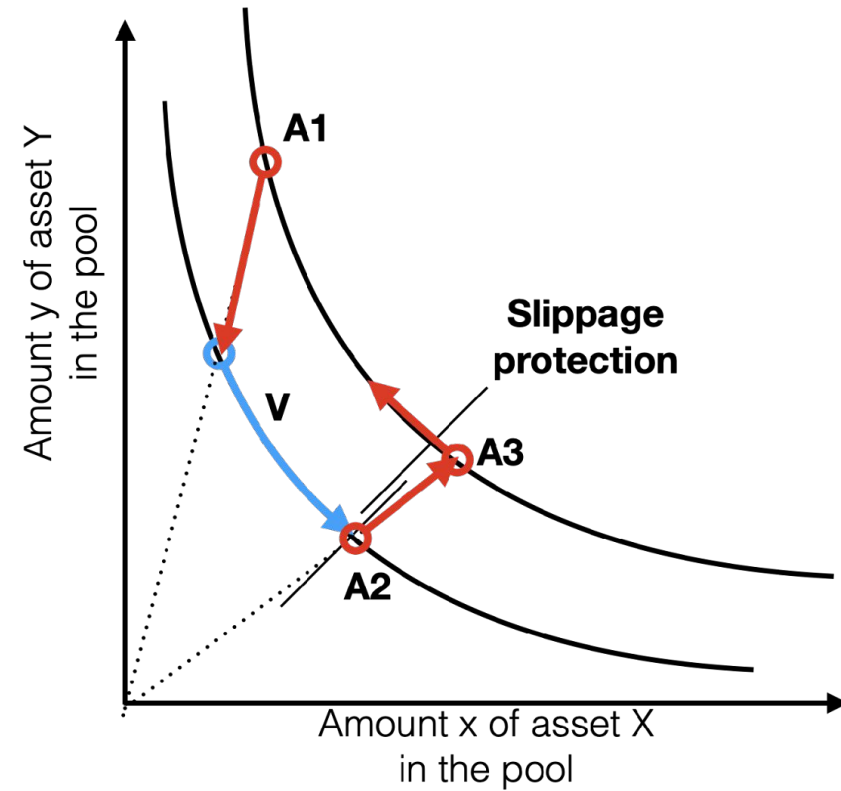


# Advanced Sandwich Attack

Taker attacks Taker



Provider attacks Taker





# Blockchain Extractable Value

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# What is Blockchain (or Miner) Extractable Value?

Price of collateral drops below health factor



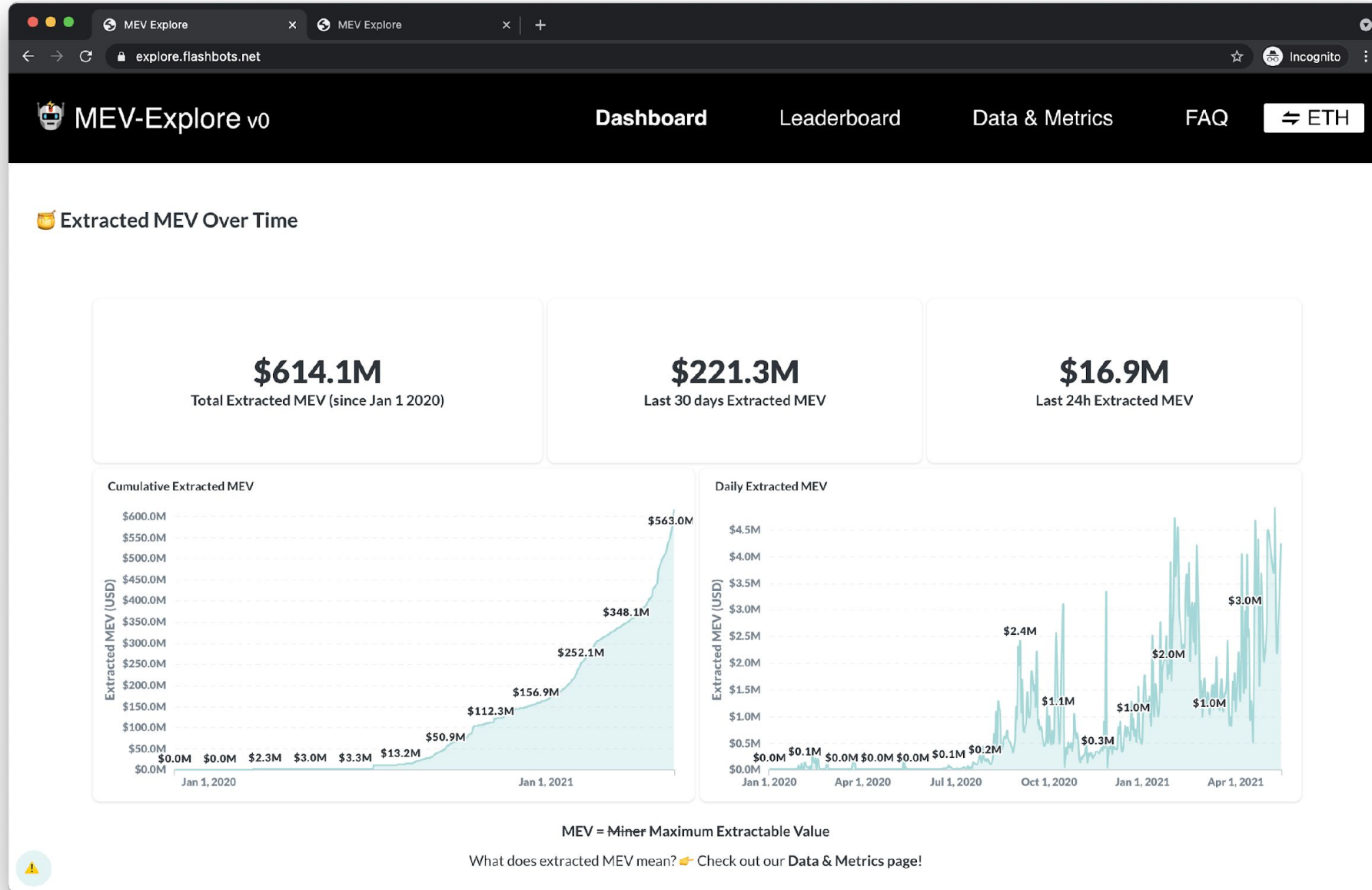
Liquidation! 



Who will liquidate?



# How much MEV?



# How much MEV? – Sandwich Attacks

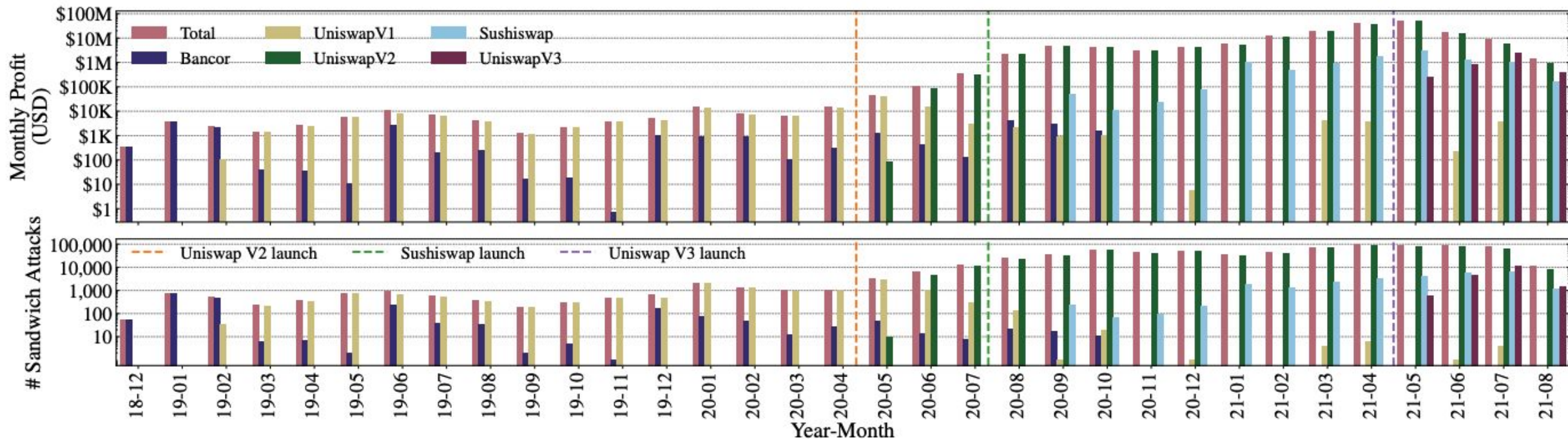
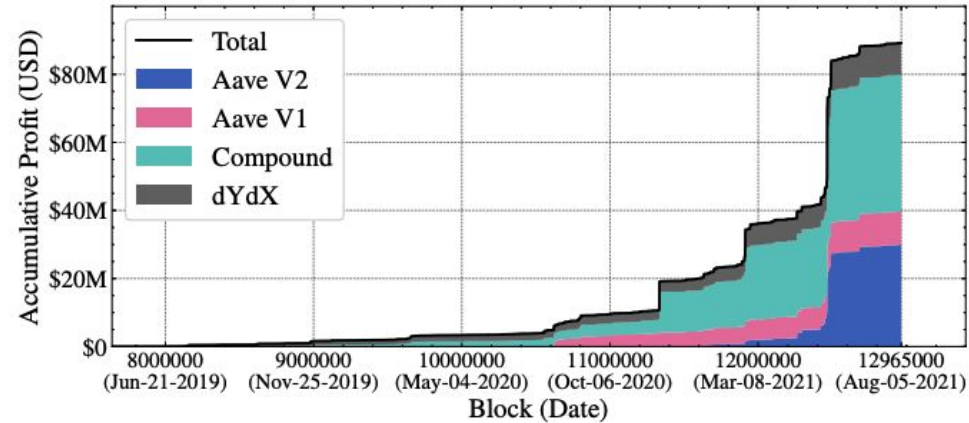
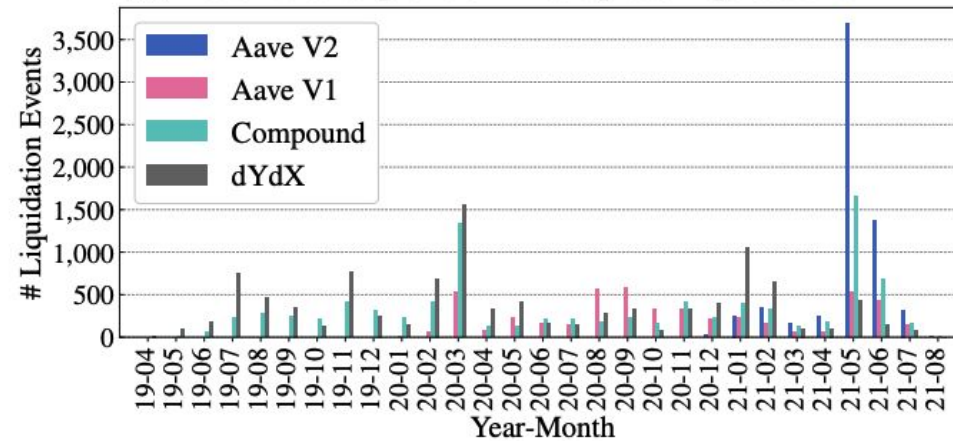


Fig. 3: Sandwich attacks, from block [6803256](#) (1st of December, 2018) to [12965000](#) (5th of August, 2021).

# How much MEV? – Liquidations



(a) Accumulative profit of fixed spread liquidations.



(b) The monthly number of fixed spread liquidation events.

Fig. 5: The number of liquidations increase in months where the ETH price collapses, e.g., in March, 2020 and May, 2021.



# How much MEV? – Arbitrage

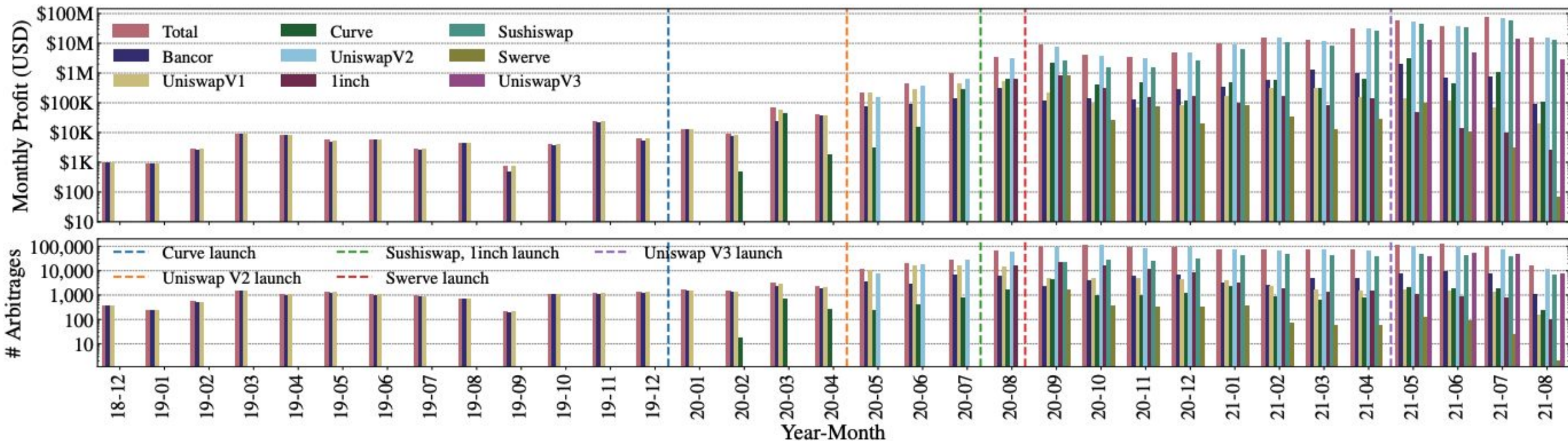


Fig. 7: Monthly arbitrage statistics from block [6803256](#) (1st of December, 2018) to block [12965000](#) (5th of August, 2021).



# Transaction Replay Attacks

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# Generalized Front-Running

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- “Copy Cat” or “Replay”
  - Observe transaction on the network layer
  - Replace certain data, sign, and broadcast copy
- Potential Profit
  - 35M USD over 32 months
  - 188,365 profitable transactions (0.02%)
  - Real-time algorithm ( $0.18s \pm 0.29$ )

# Generalized Front-Running Algorithm & Results

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## Algorithm 1: Transaction Replay Algorithm.

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**Input:** The current highest block  $B_i$ ; the potential victim transaction  $T_V$ ; the adversarial account address  $\mathcal{A}$ .

**Function** ConstructReplay( $T_V, \mathcal{A}$ ):

```

     $T.sender \leftarrow \mathcal{A}$ 
     $T.value \leftarrow T_V.value$ 
     $T.input \leftarrow$  substituting  $T_V.sender$  in  $T_V.input$  with  $\mathcal{A}$ 
    return  $T$ 
end

```

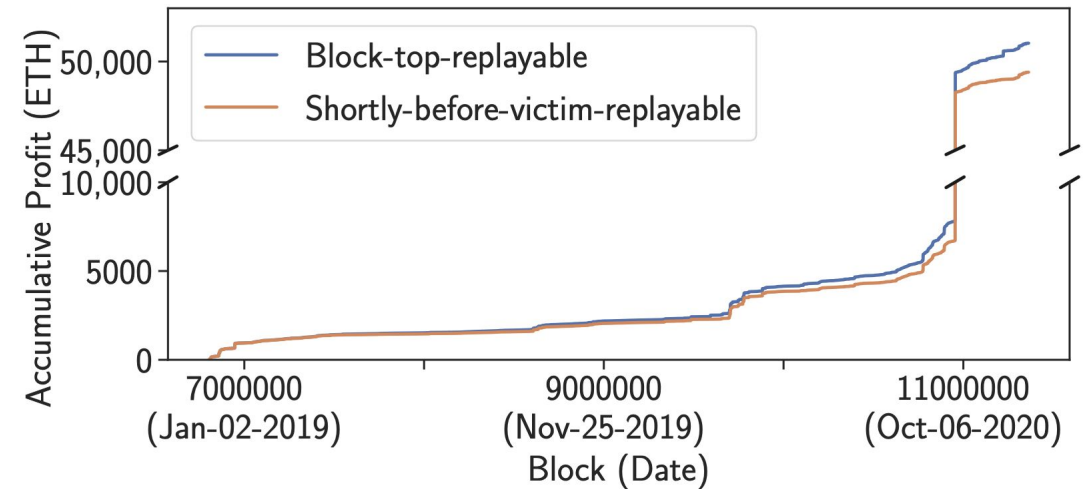
**Algorithm** TransactionReplay( $T_V, \mathcal{A}$ ):

```

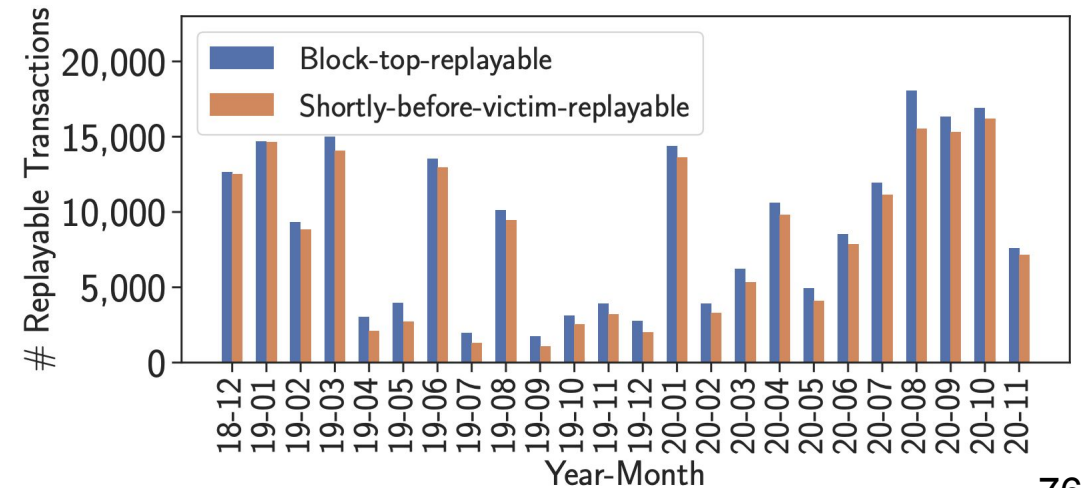
     $T_{replay} \leftarrow$  ConstructReplay( $T_V, \mathcal{A}$ )
    Concretely Execute  $T_{replay}$  upon block  $B_i$ 
    if  $T_{replay}$  is profitable then
        | Front-run  $T_V$  with  $T_{replay}$ 
    end
end

```

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(a) Accumulative profit that can be extracted by replay attacks.

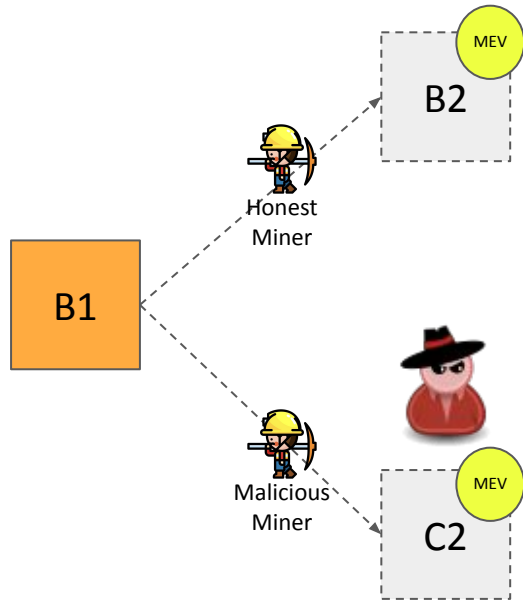




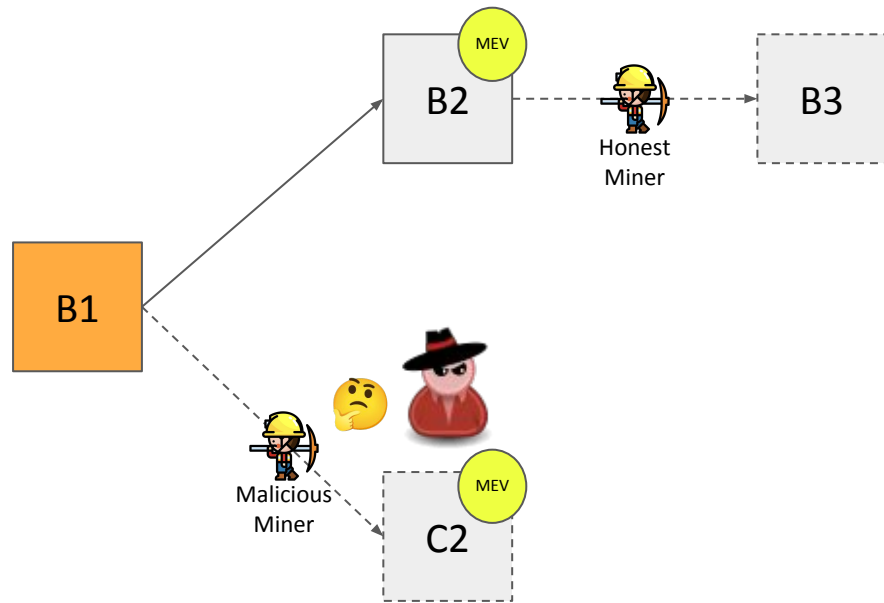
# BEV Forking and Chain Reorganisation

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# The dangers of naively maximizing MEV

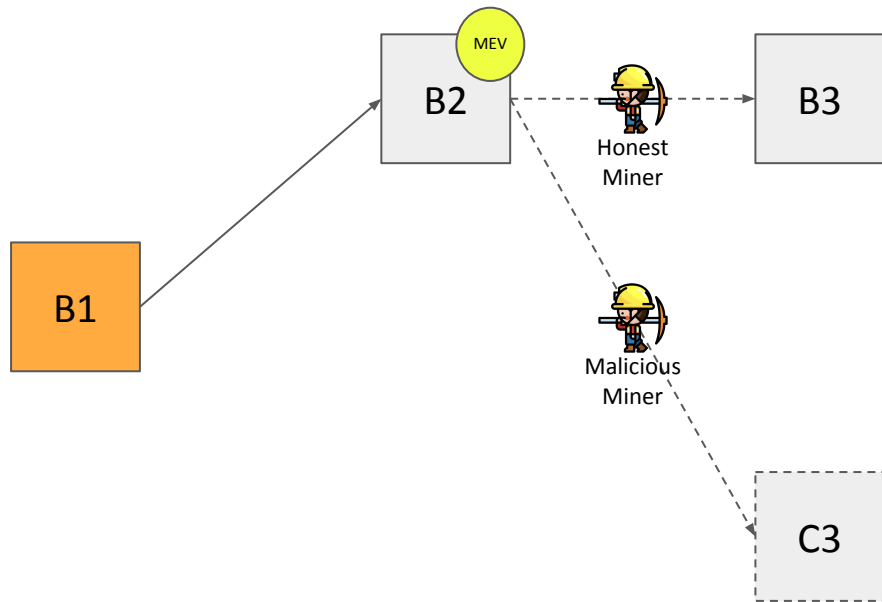


# The dangers of naively maximizing MEV



# The dangers of naively maximizing MEV

Case 1:



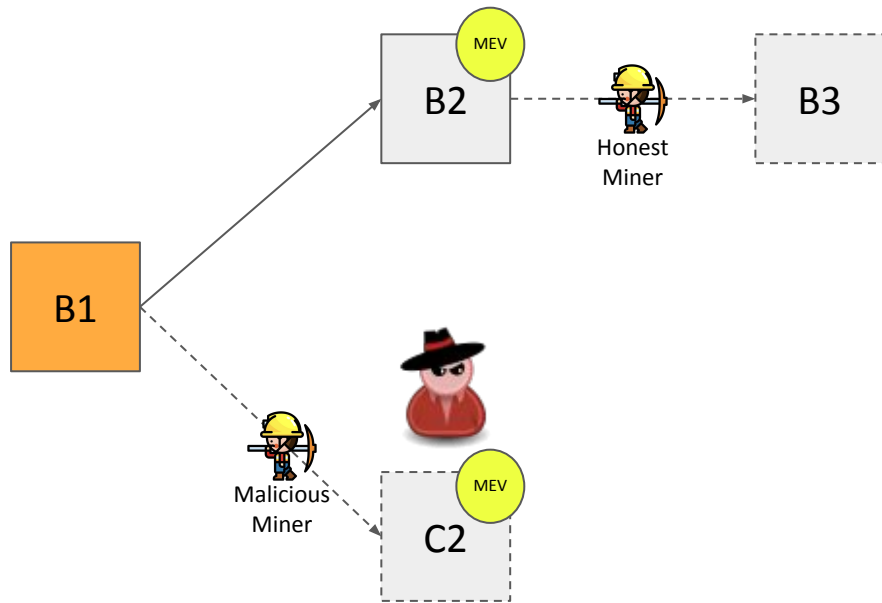
**Case 1:**

Malicious miner forfeits MEV opportunity



# The dangers of naively maximizing MEV

Case 2:



**Case 1:**

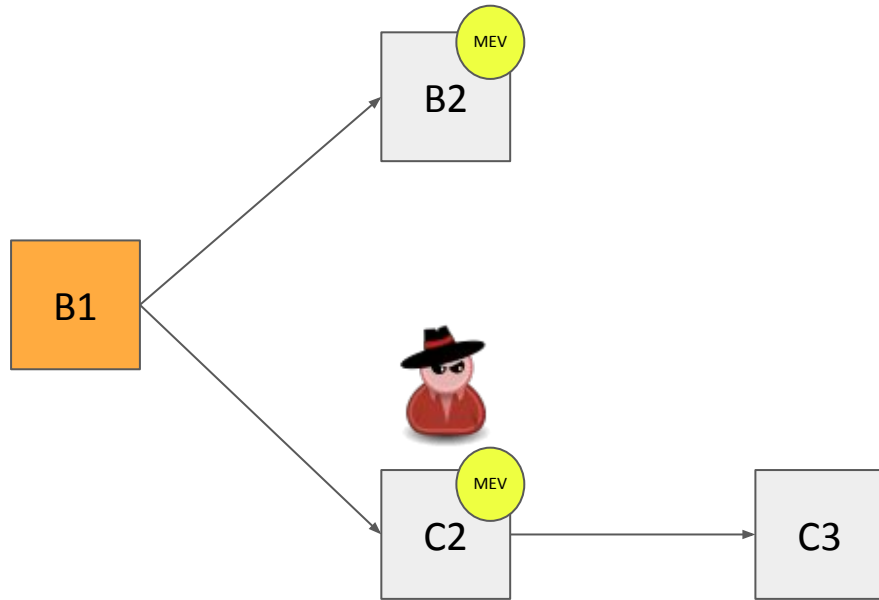
Malicious miner forfeits MEV opportunity

**Case 2:**

Keeps mining block C2

# The dangers of naively maximizing MEV

Case 2:



**Case 1:**

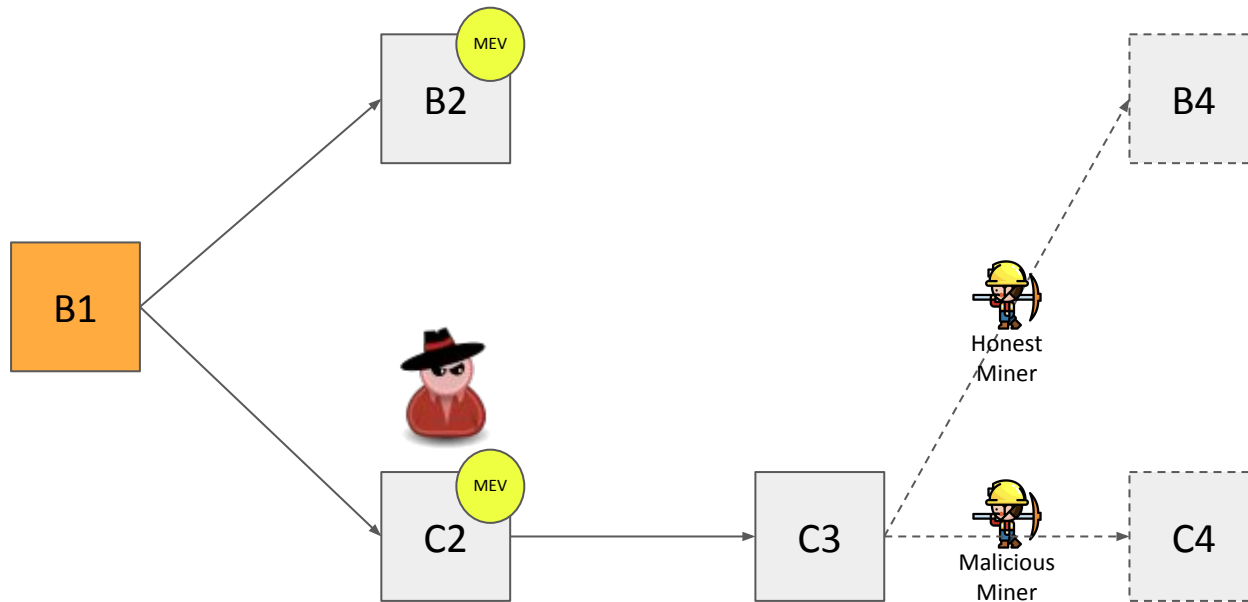
Malicious miner forfeits MEV opportunity

**Case 2:**

Keeps mining block C2

# The dangers of naively maximizing MEV

Case 2:



**Case 1:**

Malicious miner forfeits MEV opportunity

**Case 2:**

Keeps mining on block C2

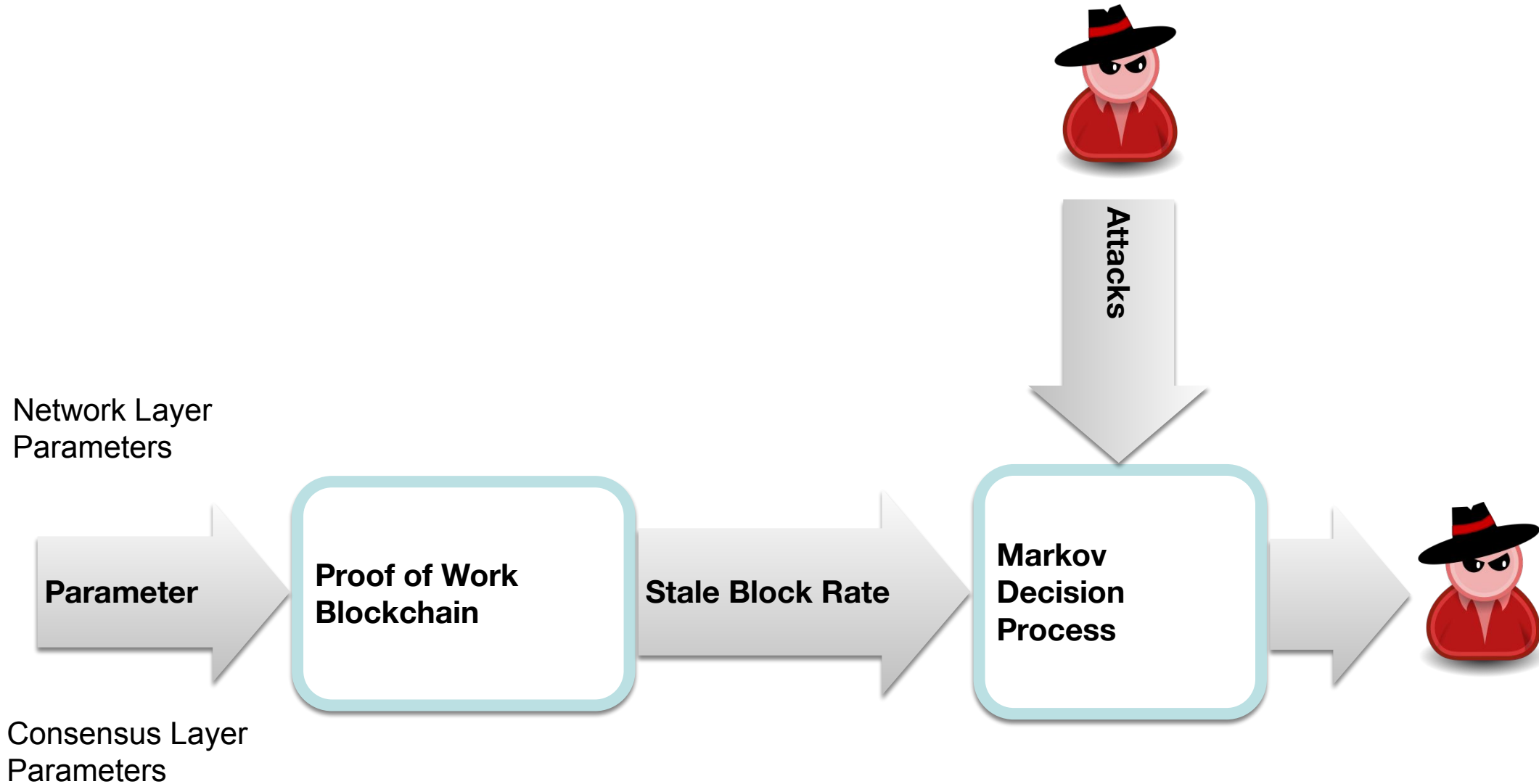
→ **Waste computational power**

→ **Increase stale block rates and risks for:**

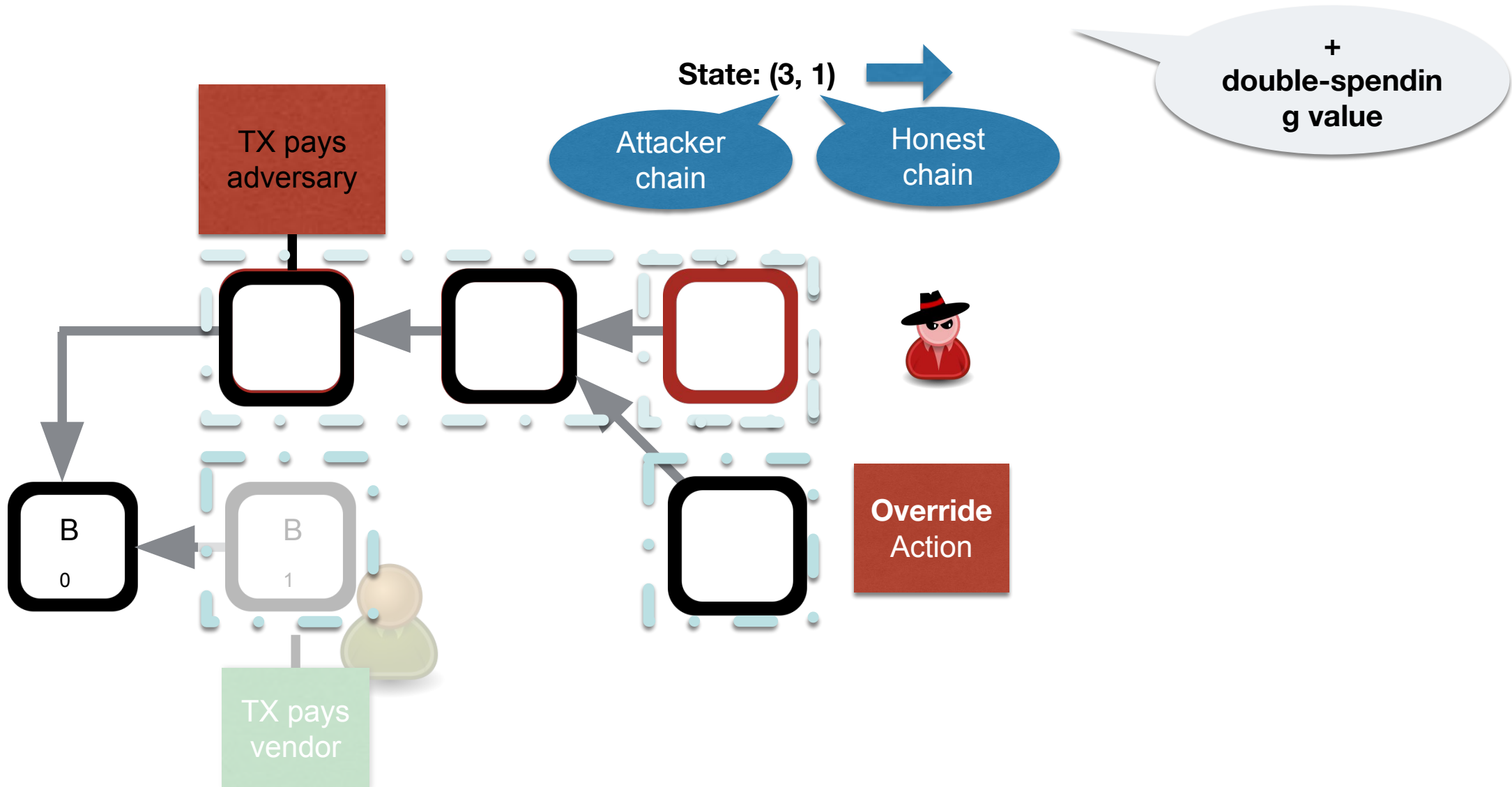
- **Double spending**
- **Selfish mining**



# Markov Decision Process (MDP)



# Markov Decision Process (MDP)



# Reducing MEV is the key to security (example)



10%  
miner

+



MEV, 4x average  
block reward

==



# Reducing MEV is the key to security (example)

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874x

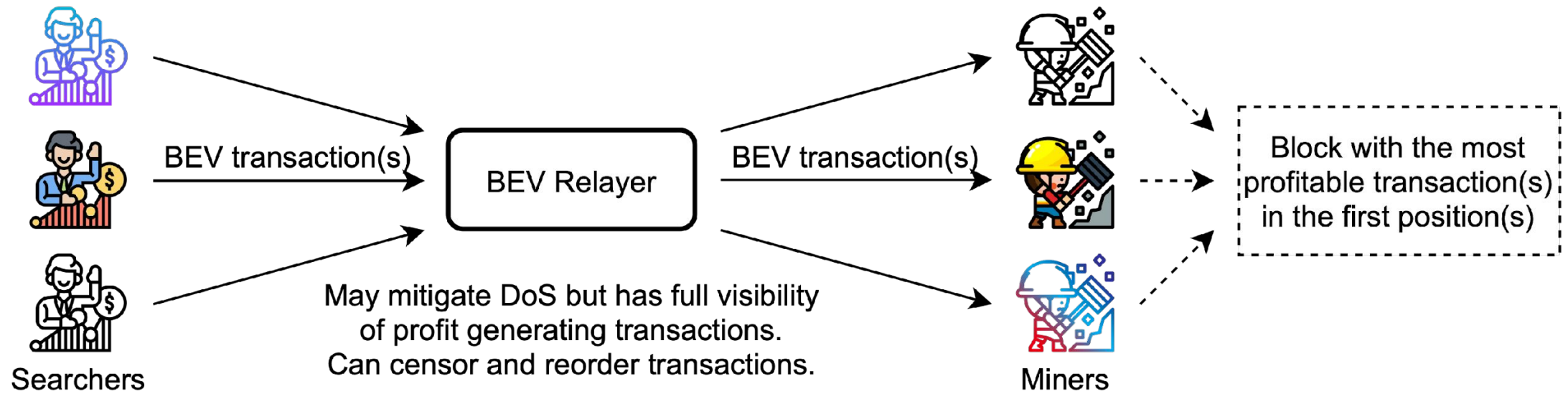


# BEV Relay & How to Mitigate BEV?

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# BEV Relay Architecture



# BEV Relay Concerns

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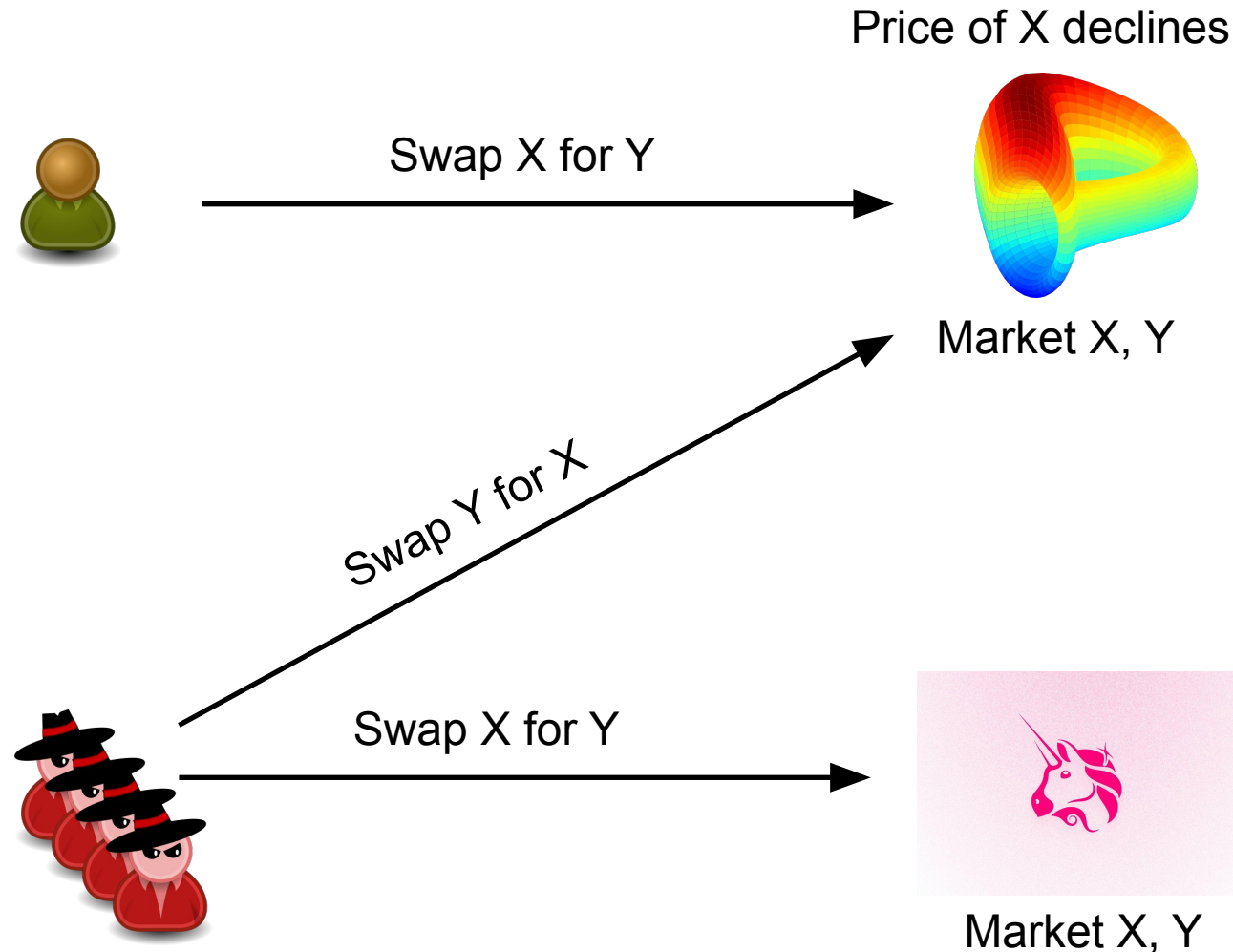
- BEV provably incentivises miners to fork (cf. S&P'21)
- BEV relayer centralise the P2P Network
- The relayer may resell/profit from searcher strategies
- The relay system doesn't necessarily reduce P2P overhead
- A for profit company distributes the geth client to >50% of the miners
- Innocent users are being stolen from systematically

# Anti-MEV Solution Space

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- Fair-Ordering on the Blockchain Layer
  - e.g., Aequitas Protocol Family
- Fixing MEV of existing dApps
  - Merging AMM DEX into one
    - On-chain aggregators such as A2MM (see DEX lecture)
- Designing MEV-Mindful dApps
  - Avoiding MEV by design
    - e.g., a price oracle update immediate performs a liquidation
- Might not fix cross-chain MEV..

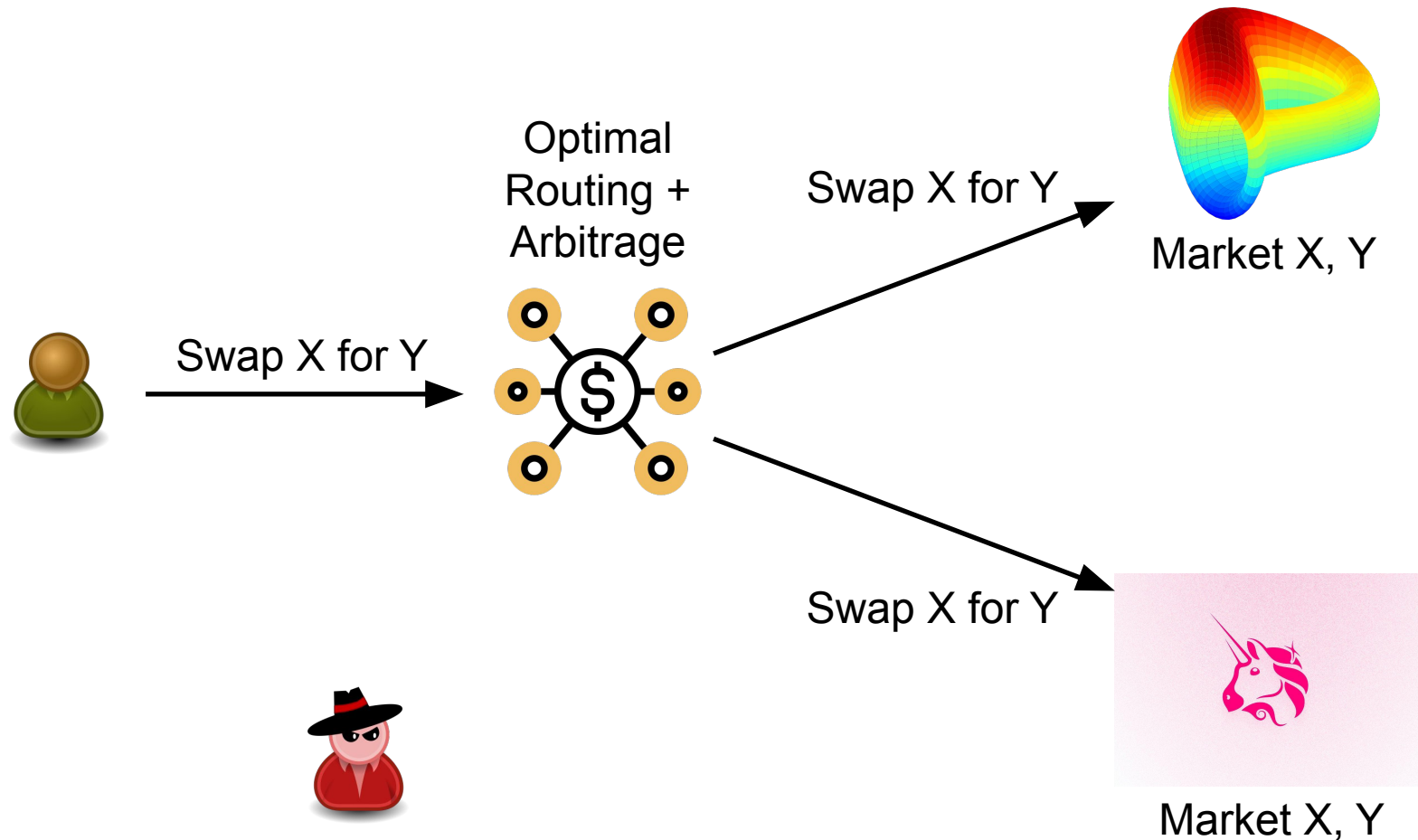
# Application-Specific MEV Mitigation



## ■ Causes

- Back-run Flooding
  - Network Congestions
  - Price Gas Auctions
  - Transaction Fee Increase
- ## ■ The user forgoes an arbitrage opportunity.

# Application-Specific MEV Mitigation



## ■ Cons

- Higher Gas Fees

## ■ Pros

- Better ex rate
- Arbitrage profit
- MEV reduction
- Healthier chain